IFT 6113 APPLICATIONS OF CURVATURE

http://tiny.cc/ift6113



Mikhail Bessmeltsev

Some slides from Alla Sheffer

MESH SIMPLIFICATION

Warning: Stone Age Tools (but still useful)



Image from Wikipedia

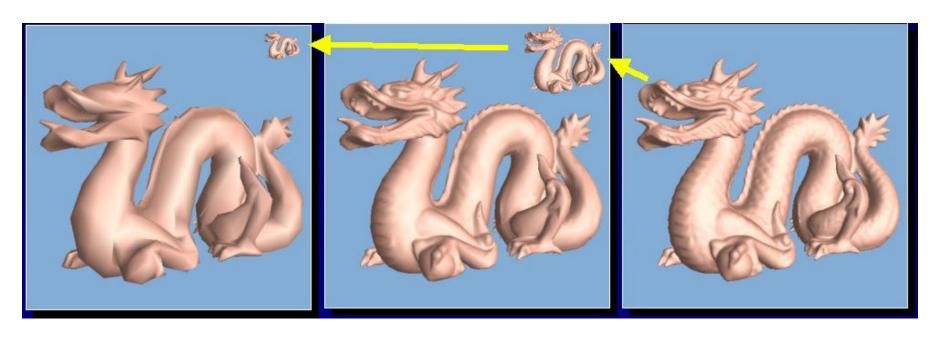


Why sim plify meshes?

- Faster rendering
- Faster collision detection
- Storage/transmission/etc.

Why sim plify meshes?

Faster rendering: Level of Detail

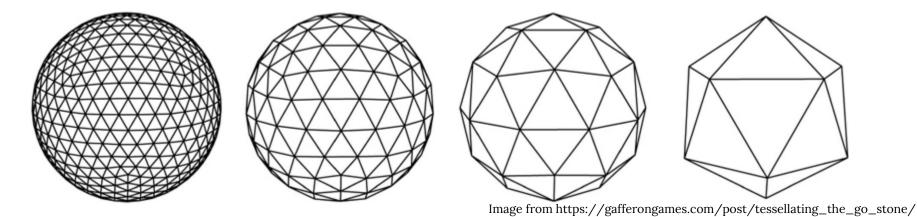


Rough, if the mesh is far

Fine, if the mesh is close

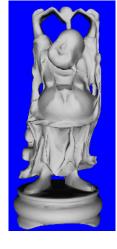
Requirements

- Reduce # of polygons
- Preserve the shape
 - Geometry
 - Features
 - Topology
 - Other constraints?

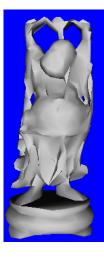


Com p le xitie s

Holes! Connected components! Boundaries!





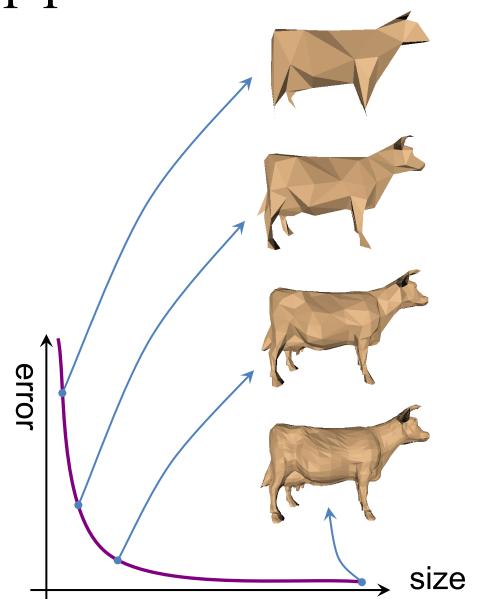


2,000



300

Approximation error



Approach: Local operations

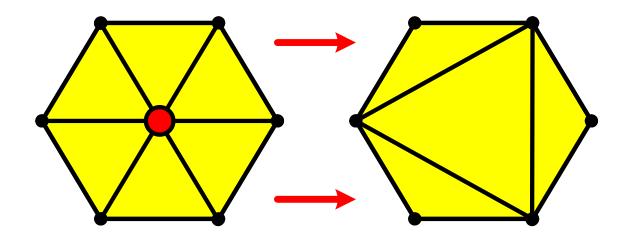
Removing edges or vertices
Keeping track of the geometry/topology

Keep score

how much will we lose if we remove a vertex?

Approach

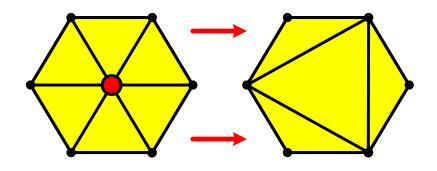
- Each operation introduces error
- Quantify?



Decim ation

Vertex removal

$$#V \rightarrow #V - 1$$
$$#F \rightarrow #F - 2$$



Remove vertex → hole → triangulate

Decim ation

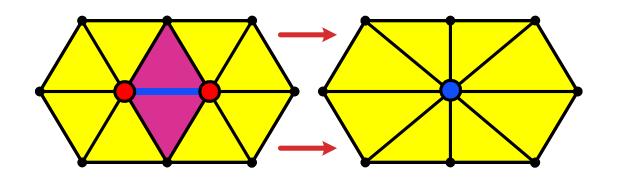
Vertex removal

$$#V \rightarrow #V - 1$$
$$#F \rightarrow #F - 2$$

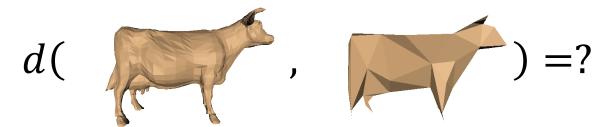
Decim ation

Edge collapse

$$#V \rightarrow #V - 1$$
$$#F \rightarrow #F - 2$$

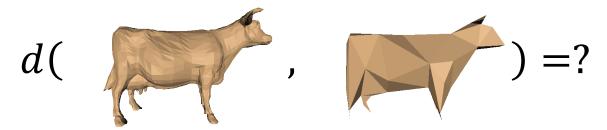


Measuring error

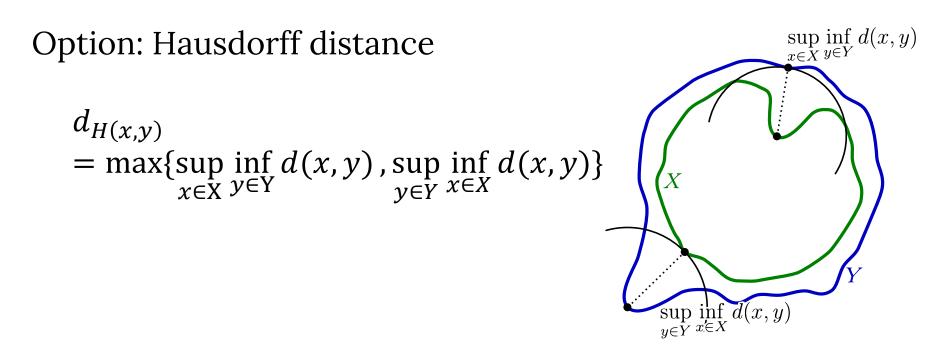


How much did we distort the shape?

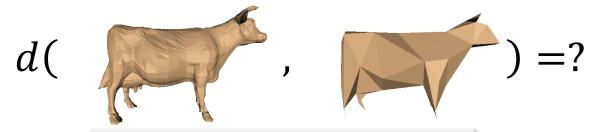
Measuring error



How much did we distort the shape?



Measuring error



How

Hard to

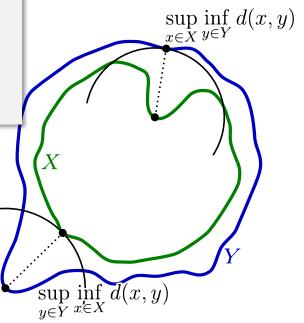
compute!

 $d_{H(x,y)}$

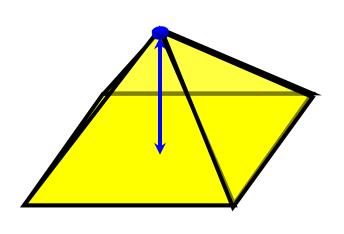
Option: Hausdo

 $= \max\{\sup_{x \in X} \inf_{y \in Y} d(x, y), \sup_{y \in Y} \inf_{x \in X} d(x, y)\}\$

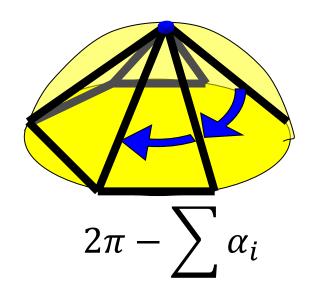
ie shape?



Local Error Control



Vertex-plane distance



Curvature

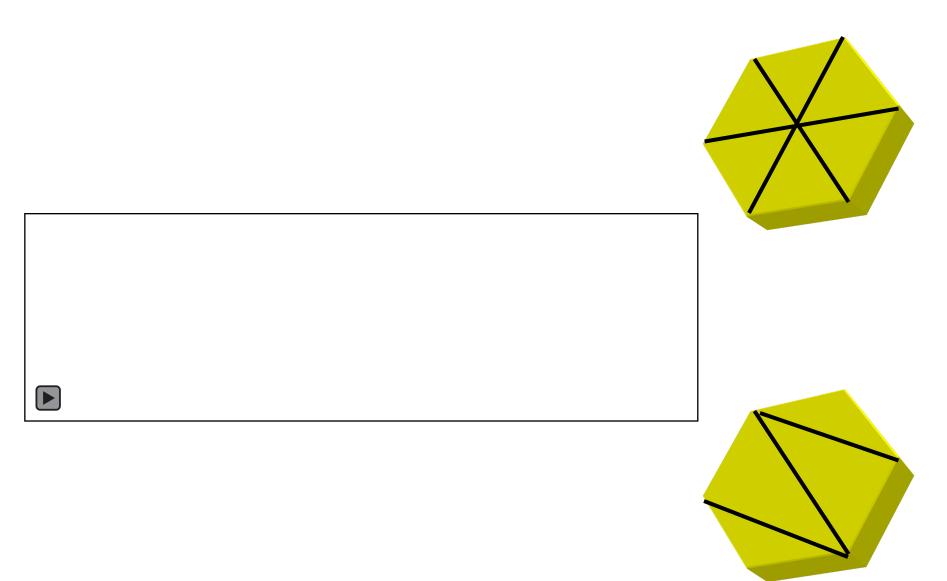
Algorithm

```
while (nVertices > nTarget)
  e = elementWithMinError()
  process(e) # remove/contract
  preserveStructure()
  updateErrors() # local/global
```

Data Structures

- Easy access to neighbors
 - Filling holes
 - Computing cost
- Priority queue/heap
 - Quick access to the cheapest element

Vertex Removal Algorithm



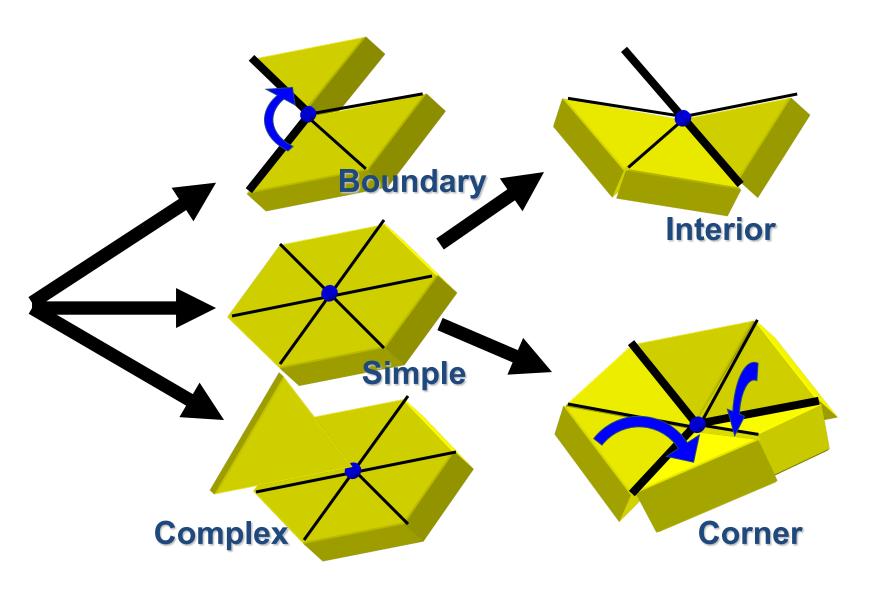
Vertex Removal Algorithm

Characterize local topology/geometry

Classify vertices as removable or not

```
while (nVertices > nTarget)
v = vertexWithMinDistance()
remove(v)
triangulateHole()
updateDistances() #to average plane
```

Characterizing structure



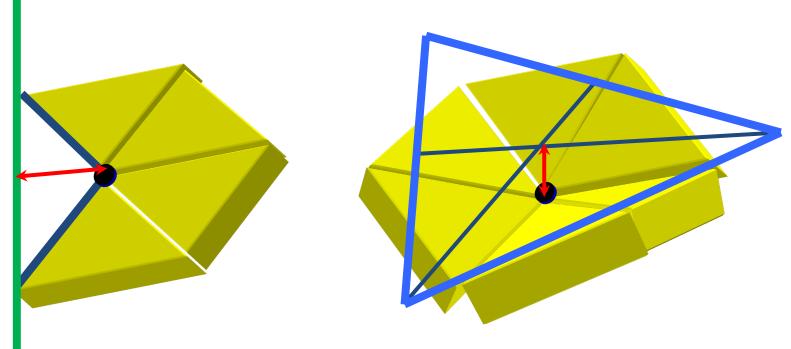
Decimation Criterion

Simple vertex *v*:

 $d(v, face loop average plane) < E_{MAX}$

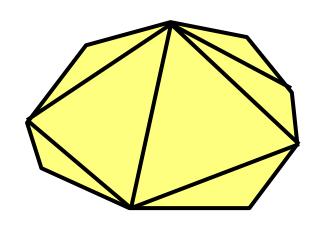
Boundary vertex v_b :

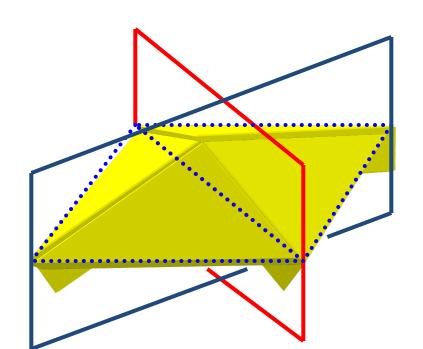
 $d(v_b, new boundary edge) < E_{MAX}$



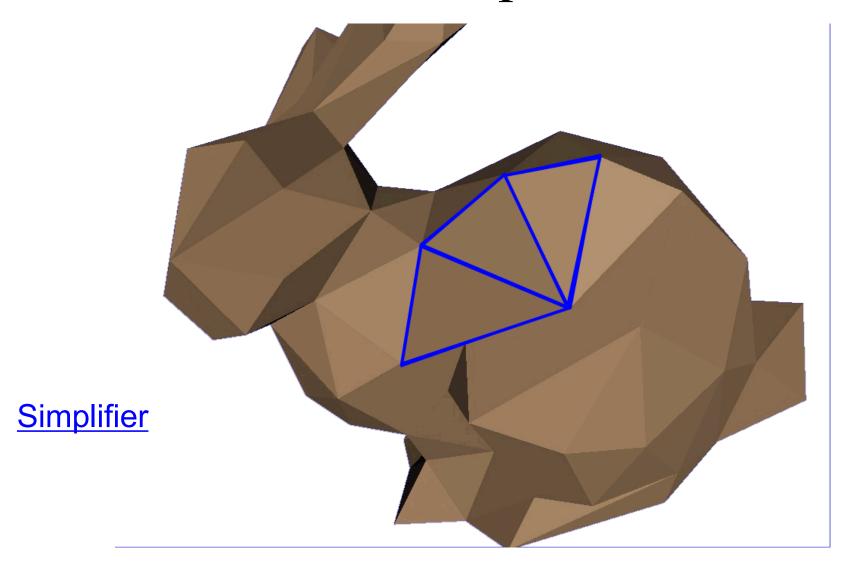
triangulateHole()

- Non-planar hole!
 - Split loop recursively
 - Split plane orthogonal to the average plane
- Control aspect ratio
- May fail
 - Vertex is not removed





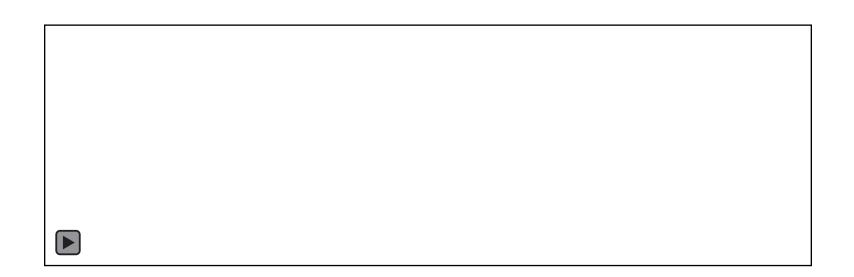
Exam ple

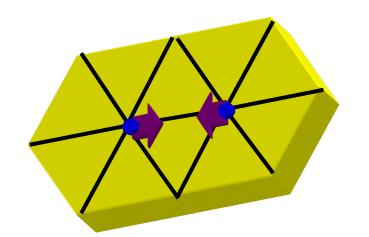


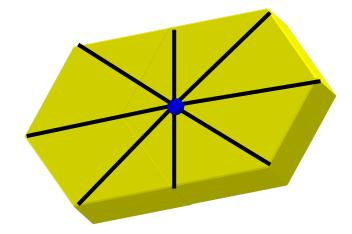
Pros and Cons

- Pros:
 - Efficient
 - Simple to implement and use
 - Few input parameters to control quality
 - Reasonable approximation
 - Works on very large meshes
 - Preserves topology
 - Vertices are a subset of the original mesh
- Cons:
 - Error is not bounded
 - Local error evaluation causes error to accumulate

Edge Collapse Algorithm

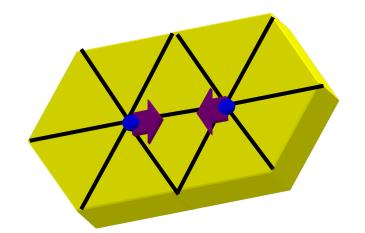


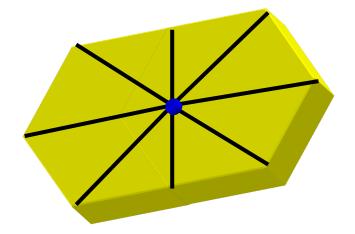




Edge Collapse Algorithm

```
while (nVertices > nTarget)
e = edgeWithMinError()
v = collapse(e) #returns a vertex
place(v)
updateErrors() # local/global
```





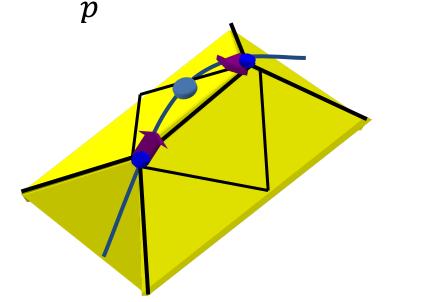
Edge Collapse Algorithm

```
while (nVertices > nTarget)
e = edgeWithMinError()
v = collapse(e) #returns a vertex
place(v)
updateErrors() # local/global
       Where should we
       place the vertex?
```

Distance Metric: Quadrics

Find the closest point to the set of adjacent triangles planes

 $v = \operatorname{argmin} d(p, planes(v))$



$$v = \underset{p}{\operatorname{argmin}} d(p, planes(v))$$

Plane equation Ax + By + Cz + D = 0

Squared distance to a plane $d(p, plane) = (Ax + By + Cz + D)^2$

assuming
$$\|\vec{n}\|^2 = A^2 + B^2 + C^2 = 1$$

Squared distance to a plane $d(p, plane) = (Ax + By + Cz + D)^2$

Squared distance to a plane

$$d(\mathbf{x}, plane) = (A\mathbf{x} + B\mathbf{y} + C\mathbf{z} + D)^{2}$$
$$= (\mathbf{x}^{T}\mathbf{p})^{2} = \mathbf{x}^{T}\mathbf{p}\mathbf{p}^{T}\mathbf{x}$$

$$\boldsymbol{p} = \begin{pmatrix} A \\ B \\ C \\ D \end{pmatrix}$$

Squared distance to a plane

$$d(\mathbf{x}, plane) = (A\mathbf{x} + B\mathbf{y} + C\mathbf{z} + D)^{2}$$
$$= (\mathbf{x}^{T}\mathbf{p})^{2} = \mathbf{x}^{T}\mathbf{p}\mathbf{p}^{T}\mathbf{x} = \mathbf{x}^{T}K\mathbf{x}$$

$$K = pp^{T} = \begin{pmatrix} A^{2} & AB & AC & AD \\ AB & B^{2} & BC & BD \\ AC & BC & C^{2} & CD \\ AD & BD & CD & D^{2} \end{pmatrix} \ge 0$$

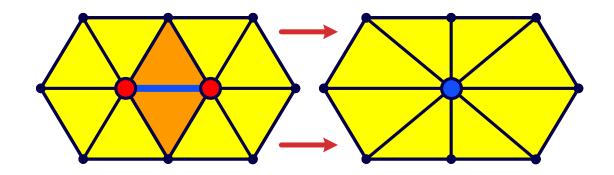
Squared distance to a set of planes

$$d(x, planes(v)) = \sum_{i} d(x, plane_{i}) =$$

$$= \sum_{i} x^{T} K_{i} x = x^{T} (\sum_{i} K_{i}) x =$$

$$= x^{T} Q_{v} x$$

$$\begin{aligned} & p \text{lace(v)} \\ & v = \mathop{\mathrm{argmin}}_{x} x^T Q_v x \\ & Q_v = Q_{v_1} + Q_{v_2} \text{\#assuming we} \\ & \text{collapsed edge between } v_1 \text{ and } v_2 \end{aligned}$$



$$v = \underset{x}{\operatorname{argmin}} x^{T} Q_{v} x$$

s. t. $x = (x, y, z, 1)^{T}$

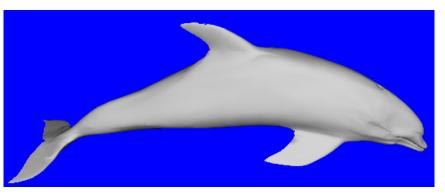
$$\Rightarrow \\ \begin{pmatrix} q_{11} & q_{12} & q_{13} & q_{14} \\ q_{12} & q_{22} & q_{23} & q_{24} \\ q_{13} & q_{23} & q_{33} & q_{34} \\ 0 & 0 & 0 & 1 \end{pmatrix} x = \begin{pmatrix} 0 \\ 0 \\ 0 \\ 1 \end{pmatrix}$$

Algorithm

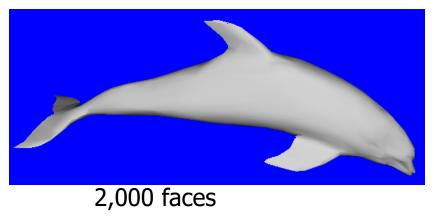
Compute Q_v for all the mesh vertices for each valid pair Compute optimal vertex position Compute its error Store all valid pairs in a priority queue (sorted by d) while reduction goal not met Take an edge from the queue, contract Update the priority queue with new valid pairs

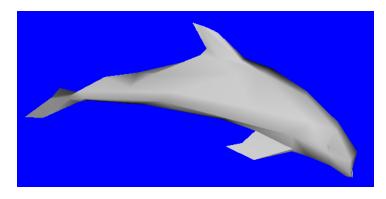
Exam p le s

Dolphin (Flipper)



Original - 12,337 faces





300 faces (142 vertices)

Pros and Cons

Pros

- Error is bounded
- Allows topology simplification
- High quality result
- Quite efficient

Cons

- Difficulties along boundaries
- Difficulties with coplanar planes
- Introduces new vertices not present in the original mesh

MESH SALIENCY

How to choose a viewpoint to show a mesh?

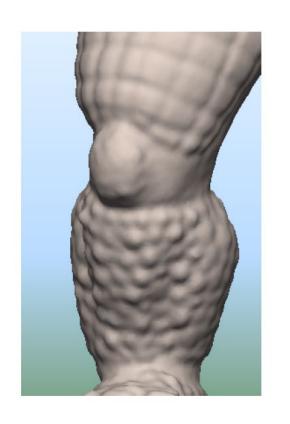


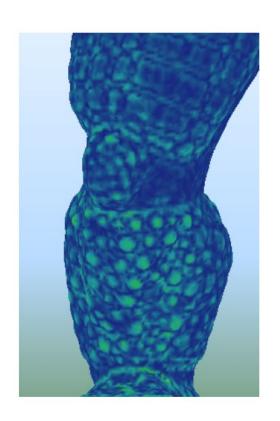


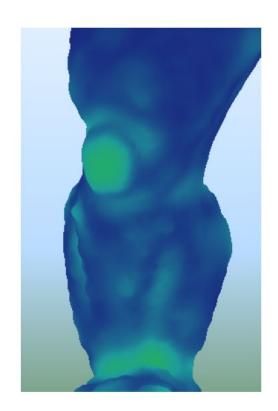
Saliency

- Gaussian curvature
- Mean curvature
- Silhouette complexity (?)

Curvature alone is not enough!







Id e a

- k_{σ} =Average mean curvature over neighborhood of radius σ
- Compute $|k_{\sigma} k_{2\sigma}|$



Original (346K triangles)



99% simplification (3.5K triangles)



98% simplification (6.9K triangles) (a) Simplification by Qslim



98.5% simplification (5.2K triangles)



99% simplification (3.5K triangles)



Saliency



99% simplification (3.5K triangles)



98% simplification 98.5 (6.9K triangles) (5 (b) Simplification guided by saliency



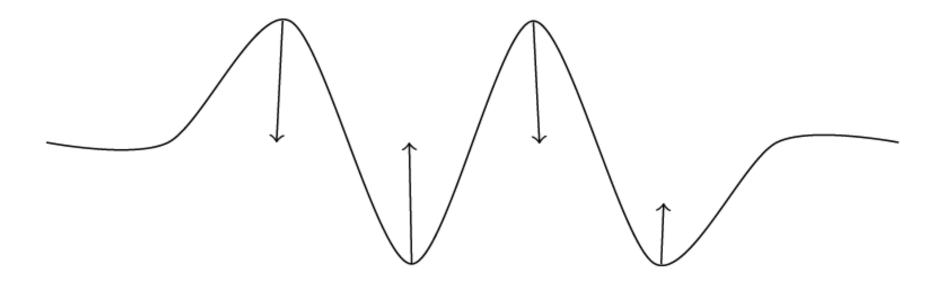
98.5% simplification (5.2K triangles)



99% simplification (3.5K triangles)

MEAN CURVATURE FLOW

What will happen to the curve?



Smoothing!

- aka Curve Shortening Flow
- (eventually) Produces convex curves!

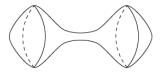


FIGURE 6. The (non-convex) dumbbell will have a "neck-pinch" before the entire surface shrinks away

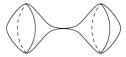


FIGURE 7. At the time of the neck-pinch, most of the surface will remain non-singular, but the diameter of the neck will have shrunk away

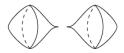
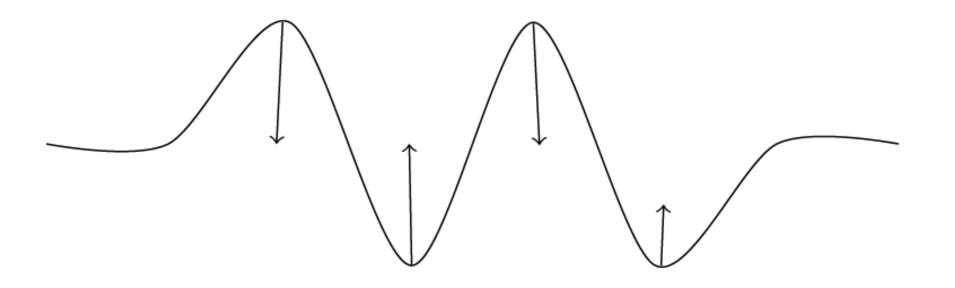


FIGURE 8. It is possible to "continue the flow through the singularity." It will become smooth immediately after, and then both component will shrink away.

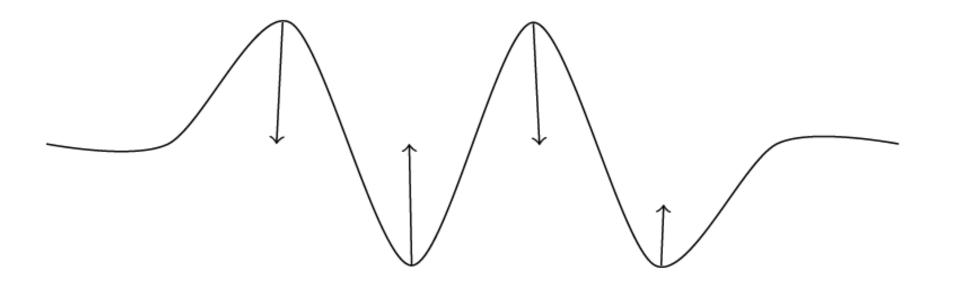
Form alization

$$\frac{\partial \gamma}{\partial T} = \frac{\partial^2 \gamma}{\partial s^2}$$



Form alization

$$\frac{\partial \gamma}{\partial T} = \frac{\partial^2 \gamma}{\partial s^2} = k\vec{n}$$



3D

Implicit Fairing of Irregular Meshes using Diffusion and Curvature Flow

Mathieu Desbrun

Mark Meyer

Peter Schröder

Alan H. Barr

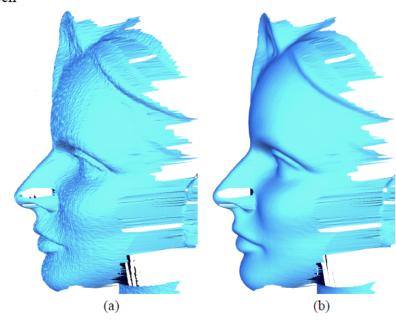
Caltech*

Abstract

In this paper, we develop methods to rapidly remove rough features from irregularly triangulated data intended to portray a smooth surface. The main task is to remove undesirable noise and uneven edges while retaining desirable geometric features. The problem arises mainly when creating high-fidelity computer graphics objects using imperfectly-measured data from the real world.

Our approach contains three novel features: an *implicit integration* method to achieve efficiency, stability, and large time-steps; a scale-dependent Laplacian operator to improve the diffusion process; and finally, a robust curvature flow operator that achieves a smoothing of the shape itself, distinct from any parameterization. Additional features of the algorithm include automatic exact volume preservation, and hard and soft constraints on the positions of the points in the mesh.

We compare our method to previous operators and related algorithms, and prove that our curvature and Laplacian operators have several mathematically-desirable qualities that improve the appearance of the control of t



Laplacian smoothing

$$\frac{\partial x_i}{\partial t} = -\overline{\mathbf{\kappa}}_i \; \mathbf{n}_i$$

$$-\overline{\kappa} \mathbf{n} = \frac{1}{4 A} \sum_{j \in N_1(i)} (\cot \alpha_j + \cot \beta_j) (x_j - x_i)$$

Laplacian Smoothing

