1. The Simplex Method

Restaurant owner problem

• Seafoods available:

30 sea-urchins24 shrimps

18 oysters

max	8x + 6y
Subject to	
	$5x + 3y \le 30$
	$2x + 3y \le 24$
	$1x + 3y \le 18$
	$x, y \ge 0$

• Two types of seafood plates to be offered: \$8 : including 5 sea-urchins, 2 shrimps et 1 oyster \$6 : including 3 sea-urchins, 3 shrimps et 3 oysters

• **Problem:** determine the number of each type of plates to be offered by the owner in order to maximize his revenue according to the seafoods available

Transforming max problem into min problem

• Consider the following maximisation problem $\max_{w} f(w)$

Subject to $w \in X \subset R^n$

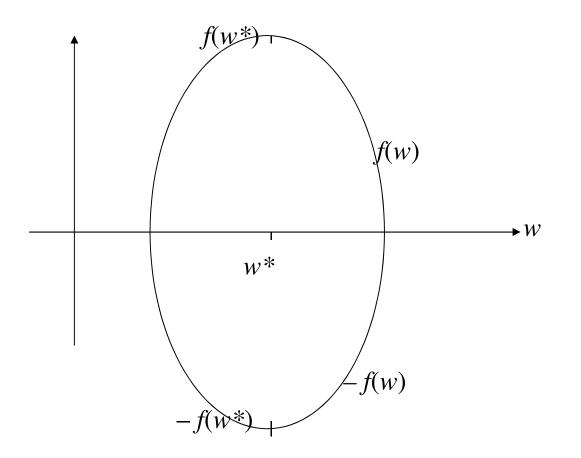
where $f: X \to R^1$.

- Let w^* be a point of X where the maximum of f is reached.
- Then $f(w^*) \ge f(w)$ $\forall w \in X$ or $-f(w^*) \le -f(w)$ $\forall w \in X$
- Consequently

 $-f(w^*) = \min - f(w)$ Subject to $w \in X \subset \mathbb{R}^n$

and w^* is a point of X where the function -f(w) reaches its minimum.

• Hence if we max f(w) or if we min -f(w) we find the same optimal solution w^* .



Transforming max problem into min problem

• Furthermore

$$f(w^*) = \max f(w) = -\min - f(w) = -(-f(w^*))$$

- We will always transform max problem into min problem.
- Then $f(w^*) \ge f(w)$ or $-f(w^*) \le -f(w)$ $\forall w \in X$ $\forall w \in X$
- Consequently

$$-f(w^*) = \min - f(w)$$

Subject to $w \in X \subset \mathbb{R}^n$

and w^* is a point of X where the function -f(w) reaches its minimum.

• Hence if we max f(w) or if we min -f(w) we find the same optimal solution w^* .

Restaurant owner problem

 $\max 8x + 6y$ Subject to $5x + 3y \le 30$ $2x + 3y \le 24$ $1x + 3y \le 18$ $x, y \ge 0$

min - (8x + 6y)Subject to $5x + 3y \le 30$ $2x + 3y \le 24$ $1x + 3y \le 18$ $x, y \ge 0$

Solving a problem graphically

- Method to solve problem having only two variables
- Consider the restaurant owner problem transformed into a min problem:

min
$$z = -8x - 6y$$

Subject to
 $5x + 3y \le 30$
 $2x + 3y \le 24$
 $1x + 3y \le 18$
 $x, y \ge 0$

 $5x + 3y \le 30$ $2x + 3y \le 24$ $1x + 3y \le 18$ $x \ge 0, y \ge 0$

• Draw the line

5x + 3y = 30

The set of points satisfying the constraint

 $5x + 3y \le 30$

is under the line since the origin satisfies the constraint

 $5x + 3y \le 30$ $2x + 3y \le 24$ $1x + 3y \le 18$ $x \ge 0, y \ge 0$

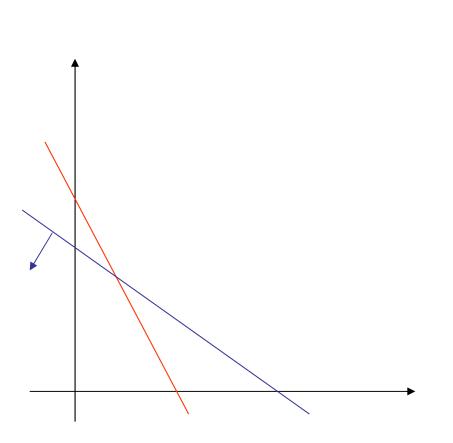
• Draw the line

2x + 3y = 24

The set of points satisfying the constraint

 $2x + 3y \le 24$

is under the line since the origin satisfies the constraint



 $5x + 3y \le 30$ $2x + 3y \le 24$ $1x + 3y \le 18$ $x \ge 0, y \ge 0$

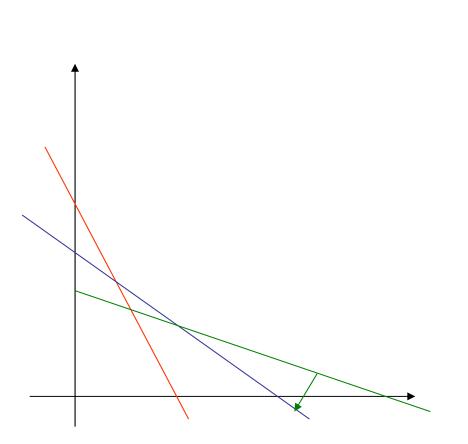
• Draw the line

x + 3y = 18

The set of points satisfying the constraint

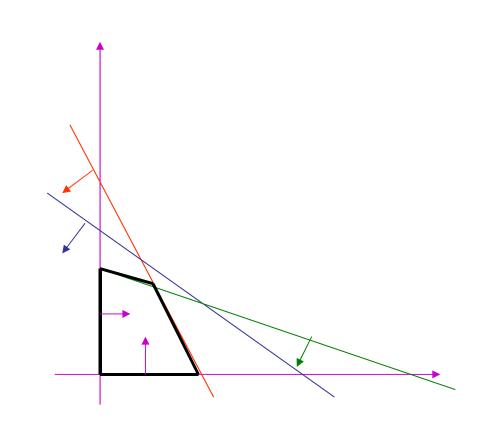
 $x + 3y \leq 18$

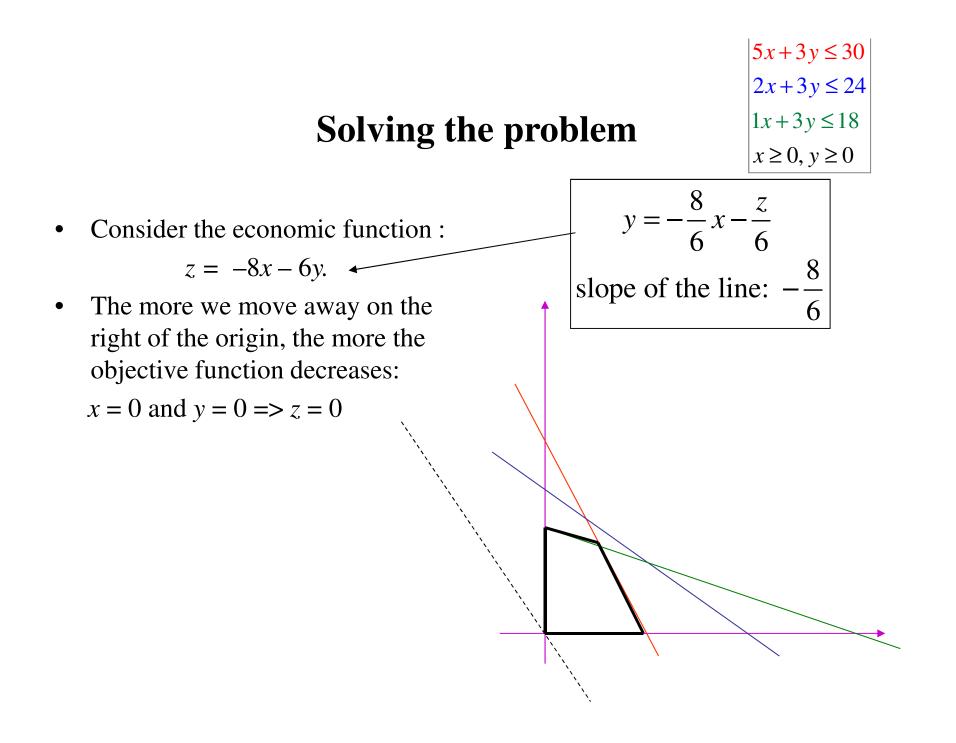
is under the line since the origin satisfies the constraint



• The set of feasible points for the system

 $5x + 3y \le 30$ $2x + 3y \le 24$ $1x + 3y \le 18$ $x, y \ge 0$





Solving the problem

 $5x+3y \le 30$ $2x+3y \le 24$ $1x+3y \le 18$ $x \ge 0, y \ge 0$

• Consider the economic function :

$$z = -8x - 6y.$$

• The more we move away on the right of the origin, the more the objective function decreases:

•
$$x = 0$$
 and $y = 0 \Longrightarrow z = 0$
 $x = 0$ and $y = 6 \Longrightarrow z = -36$

 $\begin{array}{c} x+3y=18\\ x = 0 \end{array} \Longrightarrow \begin{cases} x=0\\ y=6 \end{cases}$

Solving the problem

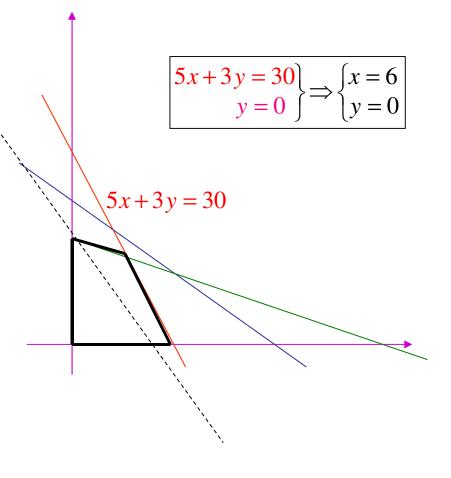
 $5x + 3y \le 30$ $2x + 3y \le 24$ $1x + 3y \le 18$ $x \ge 0, y \ge 0$

• Consider the economic function :

$$z = -8x - 6y.$$

• The more we move away on the right of the origin, the more the objective function decreases:

$$x = 0$$
 and $y = 0 \Longrightarrow z = 0$
 $x = 0$ and $y = 6 \Longrightarrow z = -36$
 $x = 6$ and $y = 0 \Longrightarrow z = -48$



Solving the problem

 $5x + 3y \le 30$ $2x + 3y \le 24$ $1x + 3y \le 18$ $x \ge 0, y \ge 0$

• Consider the economic function :

$$z = -8x - 6y.$$

• The more we move away on the right of the origin, the more the objective function decreases:

$$x = 0$$
 and $y = 0 \implies z = 0$

$$x = 0$$
 and $y = 6 \Longrightarrow z = -36$

$$x = 6$$
 and $y = 0 \Longrightarrow z = -48$

$$x = 3$$
 and $y = 5 \implies z = -54$.

• Cannot move further on the right without going out of the feasible domain.

$$\begin{bmatrix} 5x + 3y = 30\\ x + 3y = 18\\ \hline 4x = 12 \end{bmatrix} \Rightarrow \begin{cases} x = 3\\ 3 + 3y = 18 \end{cases} \Rightarrow \begin{cases} x = 3\\ y = 5 \end{cases}$$
Optimal solution:

$$x = 3 \text{ et } y = 5$$
Optimal value:

$$z = -54$$

$$x + 3y = 18$$

$$5x + 3y = 30$$

Slack variables

• Modify an inequality constraint into an equality constraint using non negative slack variables:

$$a_{i1}x_1 + a_{i2}x_2 + \dots + a_{in}x_n \le b_i \quad \to \quad a_{i1}x_1 + a_{i2}x_2 + \dots + a_{in}x_n + y_i = b_i$$

 $y_i \ge 0$

$$a_{i1}x_1 + a_{i2}x_2 + \dots + a_{in}x_n \ge b_i \longrightarrow a_{i1}x_1 + a_{i2}x_2 + \dots + a_{in}x_n - y_i = b_i$$

 $y_i \ge 0$

Minimum version of the restaurant owner

• Modify the inequality constraints of the restaurant owner problem into equality constraints using the slack variables *u*, *p*, and *h*:

min
$$z = -8x - 6y$$
min $z = -8x - 6y$ s.t. $5x + 3y \le 30$ s.t. $5x + 3y \le 30$ $5x + 3y + u$ $= 30$ $2x + 3y \le 24$ $2x + 3y + p$ $= 24$ $1x + 3y \le 18$ $1x + 3y$ $+ h = 18$ $x, y \ge 0$ $x, y, u, p, h \ge 0$

- The constraints are a linear system including 3 equations and 5 variables.
- 3 of the variables can be evaluated in terms of the other 2 variables

Dictionary Simplex Method

- min z = -8x-6ys.t. 5x + 3y + u = 30 2x + 3y + p = 24 1x + 3y + h = 18 $x, y, u, p, h \ge 0$
- The constraints are a linear system including 3 equations and 5 variables.
- 3 of the variables can be evaluated in terms of the other 2 variables

u = 30 - 5x - 3yp = 24 - 2x - 3yh = 18 - 1x - 3yz = 0 - 8x - 6y

- Fixing the values of x and y induces the values of the other 3 variables.
- It is sufficient to find non negatives values of x and y inducing non negatives values of u, p and h minimizing the value of z.
- Infinite number of possibilities. We better have a systematic procedure to find the minimum of *z*.

Find the variable to be increased

• A feasible solution of the linear system

u = 30 - 5x - 3y p = 24 - 2x - 3y h = 18 - 1x - 3y z = 0 - 8x - 6yis the following

 $x = y = 0 \implies u = 30, p = 24, h = 18 \text{ et } z = 0.$

- We can reduce the value of *z* if we increase the value of *x*, or that of *y*, or both.
- In the Simplex method we increase the value of only one
- To minimize *z*, it seems better to increase the value of *x* since increasing the value of *x* by one unit induces reducing the value of *z* by 8 units.

Limit for increasing the variable

• The non negativity of the variables u, p et h limits the increase of x

 $u = 30 - 5x - 3y \ge 0$ $p = 24 - 2x - 3y \ge 0$ $h = 18 - 1x - 3y \ge 0$

• Since the value of *y* is fixed to 0, then

 $u = 30 - 5x \ge 0 \quad \Leftrightarrow \quad x \le 30 / 5 = 6$ $p = 24 - 2x \ge 0 \quad \Leftrightarrow \quad x \le 24 / 2 = 12$ $h = 18 - 1x \ge 0 \quad \Leftrightarrow \quad x \le 18$

• The solution remains feasible as long as

 $x \le \min\{6, 12, 18\} = 6.$

New solution

•
$$u = 30 - 5x - 3y$$

 $p = 24 - 2x - 3y$
 $h = 18 - 1x - 3y$
 $z = 0 - 8x - 6y$
The solution remains feasible as long as
 $x \le \min \{6, 12, 18\} = 6.$

- In order to minimize z, we select the largest value of x:
 i.e., x = 6.
- The new solution becomes

$$x = 6, y = 0 \implies u = 0, p = 12, h = 12$$
 et $z = -48$.

New iteration

•	и		= 30 - 5x - 3y	5x + 3y + 1u	= 30
		р	= 24 - 2x - 3y	2x+3y +1p	= 24
		h	= 18 - 1x - 3y	1x + 3y + 1i	h = 18
		Z.	= 0 - 8x - 6y		

• The new solution becomes

$$x = 6, y = 0 \implies u = 0, p = 12, h = 12$$
 et $z = -48$.

- This solution is unique for the preceding system when y = u = 0 since the coefficients of the variables *x*, *p* et *h* induces a non singular matrix.
- Consequently, to determine another solution, either *y* or *u* must take a positive value.
- Previously, the analysis was simplified by the fact that the variables *x* and *y* that could be modified were on the right hand side.

- Modify the system to have *y* and *u* on the right hand side.
- Use the equation including *x* et *u* in order to find a relation where *x* is a function of *u* and *y*:

•
$$u = 30 - 5x - 3y \implies 5x = 30 - u - 3y$$

 $p = 24 - 2x - 3y$
 $h = 18 - 1x - 3y$
 $z = 0 - 8x - 6y$

- Modify the system to have *y* and *u* on the right hand side.
- Use the equation including *x* et *u* in order to find a relation where *x* is a function of *u* and *y*:
- $u = 30 5x 3y \implies (5x = 30 u 3y) \div 5$ $\implies x = 6 - \frac{1}{5u - \frac{3}{5y}}$ p = 24 - 2x - 3y h = 18 - 1x - 3yz = 0 - 8x - 6y

- Modify the system to have *y* and *u* on the right hand side.
- Use the equation including *x* et *u* in order to find a relation where *x* is a function of *u* and *y*:

•
$$u = 30 - 5x - 3y \implies x = 6 - 1/5u - 3/5y$$

 $p = 24 - 2x - 3y$
 $= p = 24 - 2(6 - 1/5u - 3/5y) - 3y$
 $= p = 12 + 2/5u - 9/5y$
 $h = 18 - 1x - 3y$
 $z = 0 - 8x - 6y$

• Replace *x* by its expression in terms of *u* and *y* in the other equations.

- Modify the system to have *y* and *u* on the right hand side.
- Use the equation including *x* et *u* in order to find a relation where *x* is a function of *u* and *y*:

•
$$u = 30 - 5x - 3y \Rightarrow x = 6 - \frac{1}{5u} - \frac{3}{5y}$$

 $p = 24 - 2x - 3y \Rightarrow p = 12 + \frac{2}{5u} - \frac{9}{5y}$
 $h = 18 - 1x - 3y$
 $= h = 18 - (6 - \frac{1}{5u} - \frac{3}{5y}) - 3y$
 $= h = 12 + \frac{1}{5u} - \frac{12}{5y}$
 $z = 0 - 8x - 6y$

• Replace *x* by its expression in terms of *u* and *y* in the other equations.

- Modify the system to have *y* and *u* on the right hand side.
- Use the equation including *x* et *u* in order to find a relation where *x* is a function of *u* and *y*:

•
$$u = 30 - 5x - 3y \Rightarrow x = 6 - 1/5u - 3/5y$$

 $p = 24 - 2x - 3y \Rightarrow p = 12 + 2/5u - 9/5y$
 $h = 18 - 1x - 3y \Rightarrow h = 12 + 1/5u - 12/5y$
 $z = 0 - 8x - 6y = z = 0 - 8(6 - 1/5u - 3/5y) - 6y$
 $=> z = -48 + 8/5u - 6/5y$

• Replace *x* by its expression in terms of *u* and *y* in the other equations.

Equivalent system

• We transformed the system

- -

_

$$= 30 - 5x - 3y \implies x \qquad = 6 - \frac{1}{5u} - \frac{3}{5y}$$

$$p \qquad = 24 - 2x - 3y \implies p \qquad = 12 + \frac{2}{5u} - \frac{9}{5y}$$

$$h \qquad = 18 - 1x - 3y \implies h \qquad = 12 + \frac{1}{5u} - \frac{12}{5y}$$

$$z = 0 \quad -8x - 6y \implies z \qquad -48 + \frac{8}{5u} - \frac{6}{5y}$$

-

Equivalent system

- We have a new system equivalent to the preceding one (i.e., the two systems have the same set of feasible solutions)
- Note that it is not interesting to increase *u* since the value of *z* would increase
- We repeat the preceding procedure by increasing the value of *y*

$$x = 6 - \frac{1}{5u} - \frac{3}{5y}$$

$$p = 12 + \frac{2}{5u} - \frac{9}{5y}$$

$$h = 12 + \frac{1}{5u} - \frac{12}{5y}$$

$$z = -48 + \frac{8}{5u} - \frac{6}{5y}$$

New iteration

• The non negativity of the variables *x*, *p* et *h* limits the increase of *y* :

$$x = 6 - \frac{1}{5u} - \frac{3}{5y} \ge 0$$

= 12 + 2/5u - 9/5y \ge 0
h = 12 + 1/5u - 12/5y \ge 0

• Since the value of *u* is fixed to 0, then

X	$= 6 - 3/5y \ge 0$	\Leftrightarrow	$y \le 10$
р	$= 12 - 9/5y \ge 0$	\Leftrightarrow	$y \le 20/3$
h	$= 12 - 12/5y \ge 0$	\Leftrightarrow	$y \le 5$

• The solution remains feasible as long as $y \le \min \{10, 20/3, 5\} = 5.$

New iteration

•
$$x = 6 - \frac{1}{5}u - \frac{3}{5}y \ge 0$$

 $p = 12 + \frac{2}{5}u - \frac{9}{5}y \ge 0$
 $h = 12 + \frac{1}{5}u - \frac{12}{5}y \ge 0$
 $z = -48 + \frac{8}{5}u - \frac{6}{5}y$
The solution remains feasible as long as
 $y \le \min \{10, \frac{20}{3}, 5\} = 5.$

- In order to minimize z, we select the largest value of y:: i.e., y = 5.
- The new solution is

$$y = 5, u = 0 \implies x = 3, p = 3, h = 0 \text{ et } z = -54.$$

Optimal solution

- Modify the system to have *h* and *u* on the right hand side.
- Use the equation including *h* and *u* in order to find a relation where *y* is a function of *h* and *u*:

$$h = 12 + 1/5u - 12/5y$$

- Replace *y* by its expression in terms of *u* and *y* in the other equations.
- The system becomes

$$x = 3 - \frac{1}{4u} + \frac{1}{4h}$$

$$p = 3 + \frac{1}{4u} + \frac{3}{4h}$$

$$y = 5 + \frac{1}{12u} - \frac{5}{12h}$$

$$z = -54 + \frac{3}{2u} + \frac{1}{2h}$$

- The solution y = 5, u = 0, x = 3, p = 3, h = 0 (where z = -54) is then optimal since the coefficients of u and h are positive.
- Indeed the value of *z* can only increase when the values of *u* or *h* increase.

It can be shown that the simplex method moves around the feasible domain to identify an optimal solution without never going inside

the feasible domain.

When solving the restaurant owner problem with the simplex method:

The initial solution is

x = y = 0 (u = 30, p = 24, h = 18) and the value of z = 0When increasing the value of x,

the solution becomes

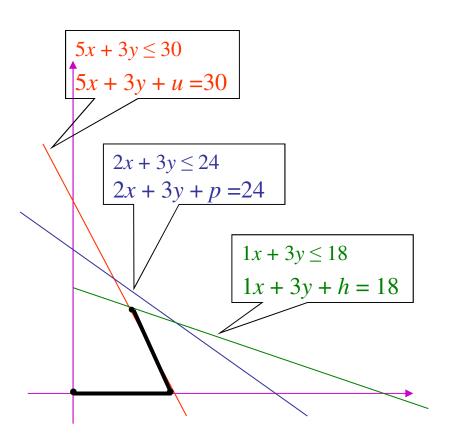
the solution becomes

x = 6, y = 0 (u = 0, p = 12, h = 12) and the value of z = -48

When increasing the value of *y*,

the solution becomes

x = 3, y = 5(u = 0, p = 3, h = 0)and the value of z = -54



Type of solutions encountered in the simplex method

- In all the solutions encountered, only 3 variables are positive!
- Since there are 5 variables, then there exist only $\binom{5}{3} = \frac{5!}{3!2!} = 10$ different solutions of this type type.
- Can there exists a solution with more than 3 positive variables having a value for *z* better than the best solution generated by the simplex.

It can be shown that this is not the case.

Standard Form

• After modifying the inequality constraints into equality constraints using slack variables, we obtain the standard form of the problem where some variables may be slack variables:

$$\begin{array}{ll} \min & z = c_1 x_1 + c_2 x_2 + \ldots + c_n x_n \\ \text{Sujet à} & a_{11} x_1 + a_{12} x_2 + \ldots + a_{1n} x_n = b_1 \\ & a_{21} x_1 + a_{22} x_2 + \ldots + a_{2n} x_n = b_2 \\ & \ddots & \ddots & \ddots \\ & \ddots & \ddots & \ddots \\ & a_{m1} x_1 + a_{m2} x_2 + \ldots + a_{mn} x_n = b_m \\ & x_1, x_2, \ldots, x_n \ge 0 \end{array}$$

Analysis of one iteration

• To analyse an iteration of the simplex method, suppose that after completing some iterations of the procedure, the variables $x_1, x_2, ..., x_m$ are function of the other variables.

The system

The system is as follows: ٠

•

•

$$x_r + + \overline{a}_{rm+1}x_{m+1} + \dots + \overline{a}_{rs}x_s + \dots + \overline{a}_{rn}x_n = \overline{b}_r$$

$$x_{m} + \overline{a}_{mm+1} x_{m+1} + \dots + \overline{a}_{ms} x_{s} + \dots + \overline{a}_{mn} x_{n} = \overline{b}_{m}$$

$$\overline{c}_{m+1} x_{m+1} + \dots + \overline{c}_{s} x_{s} + \dots + \overline{c}_{n} x_{n} = z - \overline{z}$$

•

The variables $x_1, x_2, ..., x_m$ are dependent variables of the values of the • other variables that are the independent variables.

• •

- The variables $x_1, x_2, ..., x_m$ are dependent variables of the values of the other variables that are the independent variables.
- At each iteration, we transform the system in order to maintain the non negativity of the right hand terms, and hence the dependent variables are non negative when the values of the independent variables are equal to 0.



The system

• The system is as follows

•

•

• • •

$$\begin{array}{rl} x_{1}+&&&&\\ x_{2}+&&&+\overline{a_{1m+1}}x_{m+1}+\ldots+\overline{a_{1s}}x_{s}+\ldots+\overline{a_{1n}}x_{n}=\overline{b}_{1}\\ &&&+\overline{a_{2m+1}}x_{m+1}+\ldots+\overline{a_{2s}}x_{s}+\ldots+\overline{a_{2n}}x_{n}=\overline{b}_{2} \end{array}$$

$$x_r + + \overline{a}_{rm+1}x_{m+1} + \dots + \overline{a}_{rs}x_s + \dots + \overline{a}_{rn}x_n = \overline{b}_r$$

•

•

٠

$$x_{m} + \bar{a}_{mm+1}x_{m+1} + \dots + \bar{a}_{ms}x_{s} + \dots + \bar{a}_{mn}x_{n} = \bar{b}_{m}$$

$$\bar{c}_{m+1}x_{m+1} + \dots + \bar{c}_{s}x_{s} + \dots + \bar{c}_{n}x_{n} = z - \bar{z}$$

The system

Move the independent variables on the right hand side: •

٠

$$x_{1} = \overline{b}_{1} - \overline{a}_{1m+1} x_{m+1} - \dots - \overline{a}_{1s} x_{s} - \dots - \overline{a}_{1n} x_{n}$$

$$x_{2} = \overline{b}_{2} - \overline{a}_{2m+1} x_{m+1} - \dots - \overline{a}_{2s} x_{s} - \dots - \overline{a}_{2n} x_{n}$$

$$\vdots \qquad \vdots \qquad \vdots \qquad \vdots \qquad \vdots \qquad \vdots \qquad \vdots$$

$$x_r = \overline{b}_r - \overline{a}_{rm+1} x_{m+1} - \dots - \overline{a}_{rs} x_s - \dots - \overline{a}_{rn} x_n$$

$$x_{m} = \overline{b}_{m} - \overline{a}_{mm+1} x_{m+1} - \dots - \overline{a}_{ms} x_{s} - \dots - \overline{a}_{mn} x_{n}$$
$$z = \overline{z} + \overline{c}_{m+1} x_{m+1} + \dots + \overline{c}_{s} x_{s} + \dots + \overline{c}_{n} x_{n}$$

٠

•

To select the variable to be increased (the entering variable), we look at the • z equation

$$x_{1} = \overline{b}_{1} - \overline{a}_{1m+1}x_{m+1} - \dots - \overline{a}_{1s}x_{s} - \dots - \overline{a}_{1n}x_{n}$$

= $\overline{b}_{2} - \overline{a}_{2m+1}x_{m+1} - \dots - \overline{a}_{2s}x_{s} - \dots - \overline{a}_{2n}x_{n}$

· · · ·

$$x_r = \overline{b}_r - \overline{a}_{rm+1} x_{m+1} - \dots - \overline{a}_{rs} x_s - \dots - \overline{a}_{rn} x_n$$

$$x_{m} = \overline{b}_{m} - \overline{a}_{mm+1} x_{m+1} - \dots - \overline{a}_{ms} x_{s} - \dots - \overline{a}_{mn} x_{n}$$

$$z = \overline{z} + \overline{c}_{m+1} x_{m+1} + \dots + \overline{c}_{s} x_{s} + \dots + \overline{c}_{n} x_{n}$$

•

•

• To select the variable to be increased (the entering variable), we look at the z equation

$$x_{1} = \overline{b}_{1} - \overline{a}_{1m+1}x_{m+1} - \dots - \overline{a}_{1s}x_{s} - \dots - \overline{a}_{1n}x_{n}$$

$$x_{2} = \overline{b}_{2} - \overline{a}_{2m+1}x_{m+1} - \dots - \overline{a}_{2s}x_{s} - \dots - \overline{a}_{2n}x_{n}$$

$$\vdots = \overline{b}_{r} - \overline{a}_{rm+1}x_{m+1} - \dots - \overline{a}_{rs}x_{s} - \dots - \overline{a}_{rn}x_{n}$$

$$\vdots = \overline{b}_{m} - \overline{a}_{mm+1}x_{m+1} - \dots - \overline{a}_{ms}x_{s} - \dots - \overline{a}_{mn}x_{n}$$

$$z = \overline{z} + \overline{c}_{m+1}x_{m+1} + \dots + \overline{c}_{s}x_{s} + \dots + \overline{c}_{n}x_{n}$$

• Denote
$$\overline{c}_s = \min_{1 \le j \le n} \{\overline{c}_j\}$$

• To select the variable to be increased (the entering variable), we look at the z equation

$$x_{1} = \overline{b}_{1} - \overline{a}_{1m+1}x_{m+1} - \dots - \overline{a}_{1s}x_{s} - \dots - \overline{a}_{1n}x_{n}$$

$$x_{2} = \overline{b}_{2} - \overline{a}_{2m+1}x_{m+1} - \dots$$

$$x_{r} = \overline{b}_{r} - \overline{a}_{rm+1}x_{m+1} - \dots$$

$$x_{m} = \overline{b}_{m} - \overline{a}_{mm+1}x_{m+1} - \dots$$

$$x_{m} = \overline{b}_{m} - \overline{a}_{mm+1}x_{m+1} - \dots$$

$$x_{m} = \overline{b}_{m} - \overline{a}_{mm+1}x_{m+1} + \dots + \overline{c}_{s}x_{s} + \dots + \overline{c}_{n}x_{n}$$

• Denote

$$\overline{c}_s = \min_{1 \le j \le n} \left\{ \overline{c}_j \right\}$$

• To select the variable to be increased (the entering variable), we look at the z equation

$$x_{1} = \overline{b}_{1} - \overline{a}_{1m+1}x_{m+1} - \dots - \overline{a}_{1s}x_{s} - \dots - \overline{a}_{1n}x_{n}$$

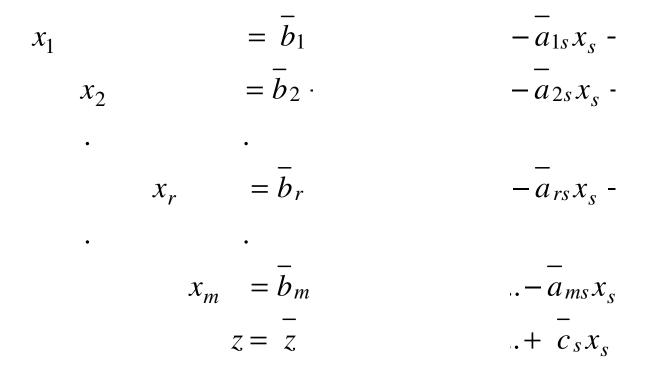
$$x_{2} = \overline{b}_{2} - \overline{a}_{2m+1}x_{m+1} - \dots - \overline{a}_{2s}x_{s} - \dots - \overline{a}_{2n}x_{n}$$
If $\overline{c}_{s} < 0$, then the variable
 x_{s} becomes the entering varaiable.
We move to Step 2.

$$x_{m} = \overline{c}_{m} + \overline{c}_{m+1}x_{m+1} + \dots + \overline{c}_{s}x_{s} - \dots - \overline{a}_{mn}x_{n}$$

$$z = \overline{z} + \overline{c}_{m+1}x_{m+1} + \dots + \overline{c}_{s}x_{s} + \dots + \overline{c}_{n}x_{n}$$

• Denote
$$\overline{c}_s = \min_{1 \le j \le n} \{\overline{c}_j\}$$

- We have to identify the largest value of the entering variable for the new solution to remain feasible.
- In fact, the increase of the entering variable is limited by the first dependent variable becoming equal to 0. This variable is denoted as the leaving variable.
- To identify the largest value of the entering variable, we refer to the preceding system of equations:



• Since the values of the other independent variables remain equal to 0, we can eliminate them from our evaluation

• The conditions insuring that the new solution remains feasible are as follows:

 $x_{1} = \overline{b}_{1} - \overline{a}_{1s} x_{s} \ge 0$ $= \overline{b}_{2} - \overline{a}_{2s} x_{s} \ge 0$ $x_{r} = \overline{b}_{r} - \overline{a}_{rs} x_{s} \ge 0$ $x_{m} = \overline{b}_{m} - \overline{a}_{ms} x_{s} \ge 0$

• Two different cases must be analysed.

• The conditions insuring that the new solution remains feasible are as follows:

 $\begin{array}{rcl} x_{1} & = \overline{b}_{1} - \overline{a}_{1s} x_{s} \geq 0 \\ x_{2} & = \overline{b}_{2} - \overline{a}_{2s} x_{s} \geq 0 \\ \vdots & \vdots & \vdots & \vdots \\ x_{r} & = \overline{b}_{r} - \overline{a}_{rs} x_{s} \geq 0 \\ \vdots & \vdots & \vdots & \vdots \\ x_{m} & = \overline{b}_{m} - \overline{a}_{ms} x_{s} \geq 0 \end{array}$ If $\overline{a}_{is} \leq 0 \quad \forall \ 1 \leq i \leq m$ then the value of the entering variable x_{s} can be increased to infinity and the values of all dependent variables remain non negative. Indeed, the value of each dependent variable x_{i} increases (if $\overline{a}_{is} < 0$) or remains the same (if $\overline{a}_{is} = 0$).

In this case, the algorithm stops indicating that the problem is not bounded below

• The conditions insuring that the new follows:

 $x_1 = \overline{b}_1 - \overline{a}_{1s} x_s \ge 0$ $= \overline{b}_2 - \overline{a}_{2s} x_s \ge 0$

. . .

. .

uring that the new In the second case, if $\overline{a}_{is} > 0$ for at least one index *i*, then the increase of the value of x_s is limited by the value when a first dependent variable reduces to 0.

$$x_r = \overline{b}_r - \overline{a}_{rs} x_s \ge 0$$

$$x_m = \overline{b}_m - \overline{a}_{ms} x_s \ge 0$$

Note that only the variables x_i where $\overline{a}_{is} > 0$ are relevant. Indeed, if $\overline{a}_{is} \le 0$, then we just observe that the value of x_i increases or remains the same, and it cannot limit the increase of the entering variable x_s .

The conditions insuring that the new solution remains feasible are as follows:
 The solution remains feasible

 $x_{1} = \overline{b}_{1} - \overline{a}_{1s} x_{s} \ge 0$ $x_{2} = \overline{b}_{2} - \overline{a}_{2s} x_{s} \ge 0$ $x_{r} = \overline{b}_{r} - \overline{a}_{rs} x_{s} \ge 0$ $x_{r} = \overline{b}_{r} - \overline{a}_{rs} x_{s} \ge 0$ $x_{m} = \overline{b}_{m} - \overline{a}_{ms} x_{s} \ge 0$ $x_{m} = \overline{b}_{m} - \overline{a}_{ms} x_{s} \ge 0$

The conditions insuring that the new solution remains feasible are as • follows:

•

 x_2

 x_1

$$x_r = \overline{b}_r - \overline{a}_{rs} x_s \ge 0$$

. .

$$x_m = \overline{b}_m - \overline{a}_{ms} x_s \ge 0$$

The solution remains feasible $= \overline{b}_{1} - \overline{a}_{1s} x_{s} \ge 0$ = $\overline{b}_{2} - \overline{a}_{2s} x_{s} \ge 0$ $\Leftrightarrow \quad \forall i \text{ such that } \overline{a}_{is} > 0$ $x_i = \overline{b}_i - \overline{a}_{is} x_s \ge 0 \quad \Leftrightarrow \quad x_s \le \frac{\overline{b}_i}{\overline{b}_i}$ a_{is}

> Consequently, the largest value of the entering variable x_s is

$$x_{s} = \frac{\overline{b}_{r}}{\overline{a}_{rs}} = \min_{1 \le i \le m} \left\{ \frac{\overline{b}_{i}}{\overline{a}_{is}} : \overline{a}_{is} > 0 \right\}$$

The conditions insuring that the new solution remains feasible are as follows:
 The solution remains feasible

 $x_1 = \overline{b}_1 - \overline{a}_{1s} x_s \ge 0$

$$x_2 \qquad = b_2 - a_{2s} x_s \ge 0$$

$$x_r = \overline{b}_r - \overline{a}_{rs} x_s \ge 0$$

$$x_m = \overline{b}_m - \overline{a}_{ms} x_s \ge 0$$

The solution remains leasible $\Leftrightarrow \forall i \text{ such that } \overline{a}_{is} > 0$ $x_i = \overline{b}_i - \overline{a}_{is} x_s \ge 0 \iff x_s \le \frac{\overline{b}_i}{\overline{a}_{is}}$

Consequently, the largest value of the entering variable x_s is

$$x_{s} = \frac{\overline{b}_{r}}{\overline{a}_{rs}} = \min_{1 \le i \le m} \left\{ \frac{\overline{b}_{i}}{\overline{a}_{is}} : \overline{a}_{is} > 0 \right\}$$

The independent variable x_r limiting the increase of the entering variable x_s is the leaving variable.

• Transform the system :

$$z = z + c_{m+1}x_{m+1} + \dots + c_s x_s + \dots + c_n x_n$$

—

• to take the entering variable x_s on the left replacing the leaving variable x_r , and vice-versa.

- Indeed exchange the role of the variables x_s et x_r because
- \blacktriangleright the entering variable x_s (being an independent variable with a 0 value) becomes a dependent variable with a non negative value
- > The leaving variable x_r (being a dependent variable with a non negative value) becomes an independent variable with a 0 value
- The set of operations to complete the transformation is referred to as the pivot

$$x_{1} = \overline{b}_{1} - \overline{a}_{1m+1}x_{m+1} - \dots - \overline{a}_{1s}x_{s} - \dots - \overline{a}_{1n}x_{n}$$

$$x_{2} = \overline{b}_{2} - \overline{a}_{2m+1}x_{m+1} - \dots - \overline{a}_{2s}x_{s} - \dots - \overline{a}_{2n}x_{n}$$

$$x_{r} = \overline{b}_{r} - \overline{a}_{rm+1}x_{m+1} - \dots - \overline{a}_{rs}x_{s} - \dots - \overline{a}_{rm}x_{n}$$

$$x_{r} = \overline{b}_{m} - \overline{a}_{mm+1}x_{m+1} - \dots - \overline{a}_{ms}x_{s} - \dots - \overline{a}_{mn}x_{n}$$

$$x_{m} = \overline{b}_{m} - \overline{a}_{mm+1}x_{m+1} - \dots - \overline{a}_{ms}x_{s} - \dots - \overline{a}_{mn}x_{n}$$

$$z = \overline{z} + \overline{c}_{m+1}x_{m+1} + \dots + \overline{c}_{s}x_{s} + \dots + \overline{c}_{n}x_{n}$$
Use the r^{th} equation to specify x_{s} in terms of $x_{m+1}, \dots, x_{s-1}, x_{s+1}, \dots, x_{n}, x_{r}$

$$x_{1} = \overline{b}_{1} - \overline{a}_{1m+1}x_{m+1} - \dots - \overline{a}_{1s}x_{s} - \dots - \overline{a}_{1n}x_{n}$$

$$x_{2} = \overline{b}_{2} - \overline{a}_{2m+1}x_{m+1} - \dots - \overline{a}_{2s}x_{s} - \overline{a}_{2n}x_{n}$$

$$x_{r} = \overline{b}_{r} - \overline{a}_{rm+1}x_{m+1} - \dots - \overline{a}_{rs}x_{s} - \dots - \overline{a}_{rm}$$

$$x_{r} = \overline{b}_{r} - \overline{a}_{rm+1}x_{m+1} - \dots - \overline{a}_{rs}x_{s} - \dots - \overline{a}_{rm}$$

$$x_{m} = \overline{b}_{m} - \overline{a}_{mm+1}x_{m+1} - \dots - \overline{a}_{ms}x_{s} - \dots - \overline{a}_{mn}x_{n}$$

$$z = \overline{z} + \overline{c}_{m+1}x_{m+1} + \dots + \overline{c}_{s}x_{s} + \dots + \overline{c}_{n}x_{n}$$
Replace x_{s} specified in terms of $x_{m+1}, \dots, x_{s-1}, x_{s+1}, \dots, x_{n}, x_{r}$ in each of the other equations

$$x_{1} = \left(\overline{b}_{1} - \overline{a}_{1s} \frac{\overline{b}_{r}}{\overline{a}_{rs}}\right) - \left(\overline{a}_{1m+1} - \overline{a}_{1s} \frac{\overline{a}_{rm+1}}{\overline{a}_{rs}}\right) x_{m+1} - \dots - \left(-\overline{a}_{1s} - \frac{1}{\overline{a}_{rs}}\right) x_{r} - \dots - \left(\overline{a}_{1n} - \overline{a}_{1s} \frac{\overline{a}_{rn}}{\overline{a}_{rs}}\right) x_{n}$$

$$x_{1} = \overline{b}_{1} - \overline{a}_{1m+1}x_{m+1} - \dots - \overline{a}_{1s}x_{s} - \dots - \overline{a}_{1n}x_{n}$$

$$x_{2} = \overline{b}_{2} - \overline{a}_{2m+1}x_{m+1} - \dots - \overline{a}_{2s}x_{s} - \dots - \overline{a}_{2n}x_{n}$$

$$x_{r} = \overline{b}_{r} - \overline{a}_{rm+1}x_{m+1} - \dots - \overline{a}_{rs}x_{s} - \dots - x_{n}$$

$$x_{r} = \overline{b}_{r} - \overline{a}_{rm+1}x_{m+1} - \dots - \overline{a}_{rs}x_{s} - \dots - x_{n}$$

$$x_{m} = \overline{b}_{m} - \overline{a}_{mm+1}x_{m+1} - \dots - \overline{a}_{ms}x_{s} - \dots - \overline{a}_{mn}x_{n}$$

$$z = \overline{z} + \overline{c}_{m+1}x_{m+1} + \dots + \overline{c}_{s}x_{s} + \dots + \overline{c}_{n}x_{n}$$
Replace x_{s} specified in terms of $x_{m+1}, \dots, x_{s-1}, x_{s+1}, \dots, x_{n}, x_{r}$, in each of the other equations

$$x_{1} = \overline{b}_{1} - \overline{a}_{1m+1}x_{m+1} - \dots - \overline{a}_{1s}x_{s} - \dots - \overline{a}_{1n}x_{n}$$

$$= \overline{b}_{2} - \overline{a}_{2m+1}x_{m+1} - \dots - \overline{a}_{2s}x_{s} - \dots - \overline{a}_{2n}x_{n}$$

. .

. .

•

•

$$x_r = \overline{b}_r - \overline{a}_{rm+1} x_{m+1} - \dots - \overline{a}_{rs} x_s - \dots - \overline{a}_{rn} x_n \qquad x_s = \frac{\overline{b}_r}{\overline{a}_{rs}}$$



$$x_{m} = \overline{b}_{m} - \overline{a}_{mm+1} x_{m+1} - \dots - \overline{a}_{ms} x_{s} - \dots - \overline{a}_{mn} x_{n}$$

$$z = \overline{z} + \overline{c}_{m+1} x_{m+1} + \dots + \overline{c}_{s} x_{s} + \dots + \overline{c}_{n} x_{n}$$
Replace x_{s} specified in terms of $x_{m+1}, \dots, x_{s-1}, x_{s+1}, \dots, x_{n}, x_{r}$, in each of the other equations

•

٠

.

$$x_{m} = \overline{b}_{m} - \overline{a}_{mm+1} x_{m+1} - \dots - \overline{a}_{ms} x_{s} - \dots - \overline{a}_{mn} x_{n}$$

$$z = \overline{z} + \overline{c}_{m+1} x_{m+1} + \dots + \overline{c}_{s} x_{s} + \dots + \overline{c}_{n} x_{n}$$
Replace x_{s} specified in terms of $x_{m+1}, \dots, x_{s-1}, x_{s+1}, \dots, x_{n}, x_{r}$, in each of the other equations

Equivalent system for the next iteration

• The pivot generates an equivalent system having the following form

$$\begin{array}{ll} x_{1} & = \widetilde{b}_{1} - \widetilde{a}_{1m+1}x_{m+1} - \dots - \widetilde{a}_{1r}x_{r} - \dots - \widetilde{a}_{1n}x_{n} \\ x_{2} & = \widetilde{b}_{2} - \widetilde{a}_{2m+1}x_{m+1} - \dots - \widetilde{a}_{2r}x_{r} - \dots - \widetilde{a}_{2n}x_{n} \\ \vdots & \vdots & \vdots & \vdots \\ x_{s} & = \widetilde{b}_{r} - \widetilde{a}_{rm+1}x_{m+1} - \dots - \widetilde{a}_{rr}x_{r} - \dots - \widetilde{a}_{rm}x_{n} \\ \vdots & \vdots & \vdots & \vdots \\ x_{m} & = \widetilde{b}_{m} - \widetilde{a}_{mm+1}x_{m+1} - \dots - \widetilde{a}_{mr}x_{r} - \dots - \widetilde{a}_{mn}x_{n} \\ z = \widetilde{z} + \widetilde{c}_{m+1}x_{m+1} + \dots + \widetilde{c}_{r}x_{r} + \dots + \widetilde{c}_{n}x_{n} \end{array}$$

• Using this new system, we complete a new iteration.

Tableau format of the simplex method

- We now use the tableau format to complete the iterations of the simplex method.
- We illustrate one iteration of the tableau format for solving the restaurant owner problem.

Recall the problem

min z = -8x - 6ySujet à 5x + 3y + u = 30 2x + 3y + p = 24 1x + 3y + h = 18 $x, y, u, p, h \ge 0$

Equivalent tableau format for the system

$\min \ z = -8x - 6y$	min z		
Subject to	Subject to		
5x + 3y + u = 30	5x + 3y + u = 30		
2x + 3y + p = 24	2x + 3y + p = 24		
$1x + 3y \qquad \qquad + h = 18$	$1x + 3y \qquad \qquad + h = 18$		
$x, y, u, p, h \ge 0$	$-8x - 6y \qquad -z = 0$		
	$x, y, u, p, h \ge 0$		
u = 30 - 5x - 3y	dep. var. x y u p h $-z$ r.h.s.		
p = 24 - 2x - 3y	<i>u</i> 5 3 1 30		
	<i>p</i> 2 3 1 24		
h = 18 - 1x - 3y	h 1 3 1 18		
z = 0 -8x - 6y	<i>-z</i> -8 -6 1 0		

и	= 30 - 5x - 3y	
p	= 24 - 2x - 3y	
h	= 18 - 1x - 3y	
	z = 0 -8x - 6y	

$\begin{array}{c ccc} u & 5 & 3 \\ p & 2 & 3 \end{array}$	1				30
<i>p</i> 2 3					
		1			24
$h \mid 1 \mid 3$			1		18
-z -8 -6	5			1	0

Step 1: Entering criterion

$$\overline{c}_s = \min_{1 \le j \le n} \left\{ \overline{c}_j \right\}$$

Determine the entering variable by selecting the smallest element in the last row of the tableau $min \{-8, -6, 0, 0, 0\} = -8.$ *x* is then the entering variable

$$u = 30 - 5x - 3y$$

$$p = 24 - 2x - 3y$$

$$h = 18 - 1x - 3y$$

$$z = 0 - 8x - 6y$$

r.h.s. h dep. var. - *z* 1 х U pV 3 5 30 1 U 2 3 1 24 р 3 h 1 18 1 -8 -6 1 0 -Z

Step 2: leaving criterion

$$x_s = \frac{\overline{b}_r}{\overline{a}_{rs}} = \min_{1 \le i \le m} \left\{ \frac{\overline{b}_i}{\overline{a}_{is}} : \overline{a}_{is} > 0 \right\}$$

entering variable

To identify the leaving variable determine the min of the ratio right hand side terms divided by the corresponding elements in the column of the entering variable that are positive:

$$u = 30 - 5x - 3y$$

$$p = 24 - 2x - 3y$$

$$h = 18 - 1x - 3y$$

$$z = 0 - 8x - 6y$$

$$x_{s} = \frac{\bar{b}_{r}}{\bar{a}_{rs}} = \min \left\{ \frac{\bar{b}_{i}}{\bar{a}_{is}} : \bar{a}_{is} > 0 \right\}$$

$$dep. var. x y u p h - z r.h.s.$$

$$u 5 3 1 - 30$$

$$p 2 3 - 1 - 24$$

$$h 1 3 - 1 - 18$$

$$r.h.s.$$

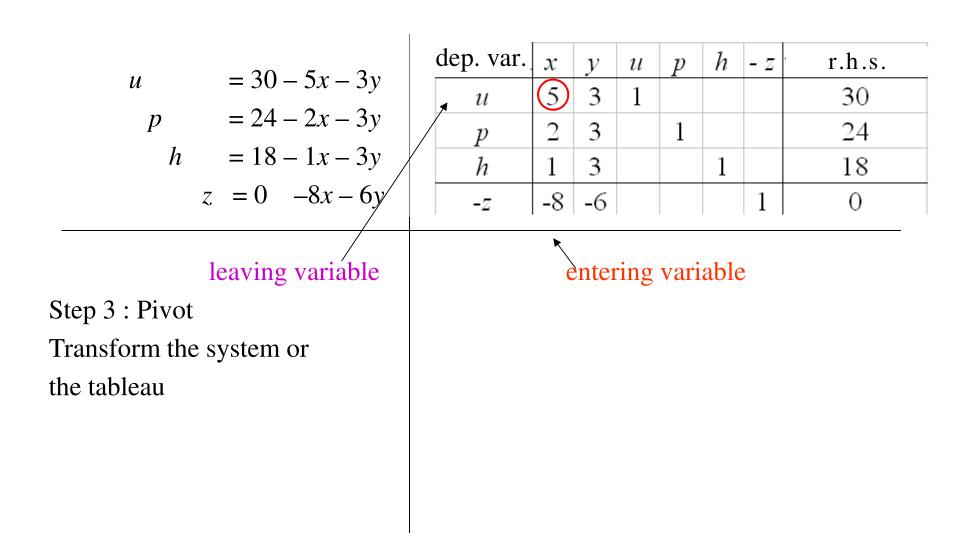
$$u -z - 8 - 6 - 1 - 1 - 0$$

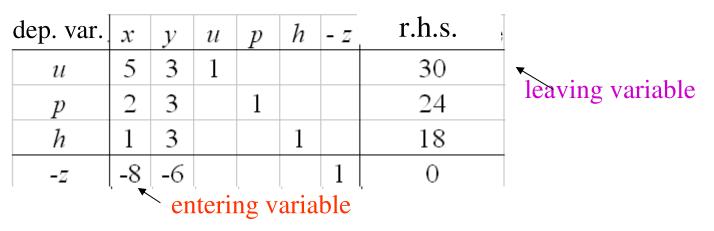
$$entering variable$$

$$min \{30/5, 24/2, 18\} = 30/5 = 6$$

$$The corresponding variable u$$

$$becomes the leaving variable$$



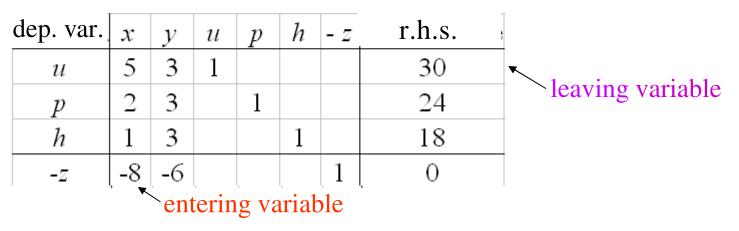


RECALL: We use the equation including variable *x* and *u* to specify *x* in terms of *u* and *y*:

$$u = 30 - 5x - 3y \implies (5x = 30 - u - 3y) / 5$$
$$\implies x = 6 - 1/5u - 3/5y$$

This is equivalent to

$$5x + 3y + u = 30$$

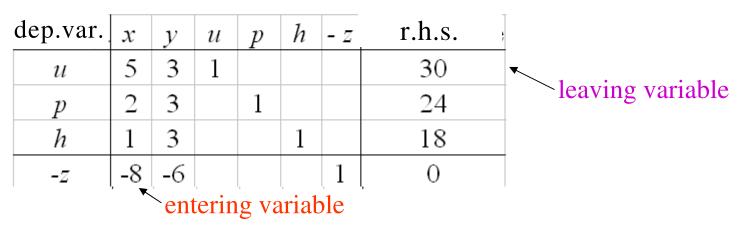


RECALL: We use the equation including variable *x* and *u* to specify *x* in terms of *u* and *y*:

$$u = 30 - 5x - 3y \implies (5x = 30 - u - 3y) / 5$$
$$\implies x = 6 - 1/5u - 3/5y$$

This is equivalent to

$$(5x + 3y + u = 30) / 5$$

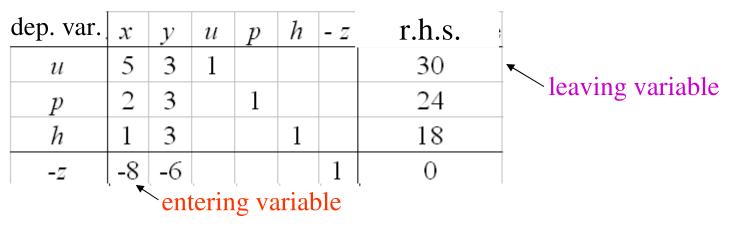


RECALL: We use the equation including variable *x* and *u* to specify *x* in terms of *u* and *y*:

$$u = 30 - 5x - 3y \implies (5x = 30 - u - 3y) / 5$$
$$\implies x = 6 - 1/5u - 3/5y$$

This is equivalent to

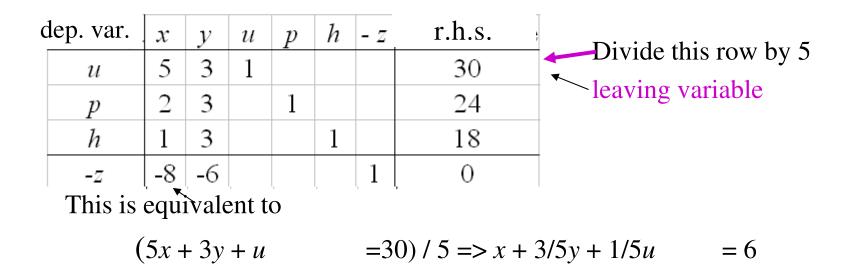
$$(5x + 3y + u = 30) / 5 => x + 3/5y + 1/5u = 6$$



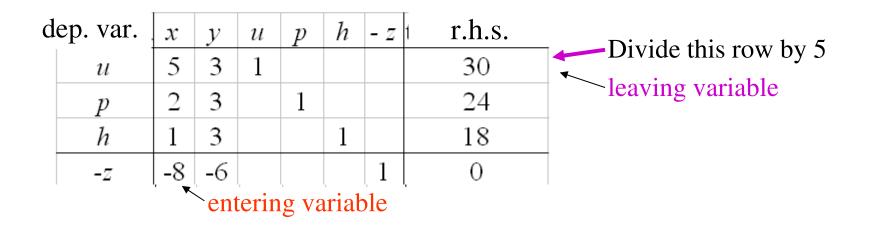
This is equivalent to

$$(5x + 3y + u = 30) / 5 => x + 3/5y + 1/5u = 6$$

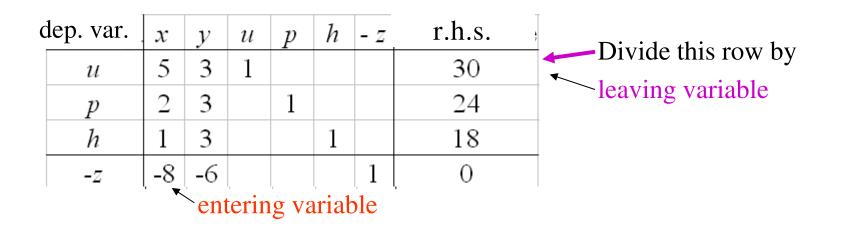
In the tableau, this is equivalent to divide the row including the leaving variable by the coefficient of the entering variable in this row



In the tableau, this is equivalent to divide the row including the leaving variable by the coefficient of the entering variable in this row



$$x + 3/5y + 1/5u = 6$$



We obtain the following tableau

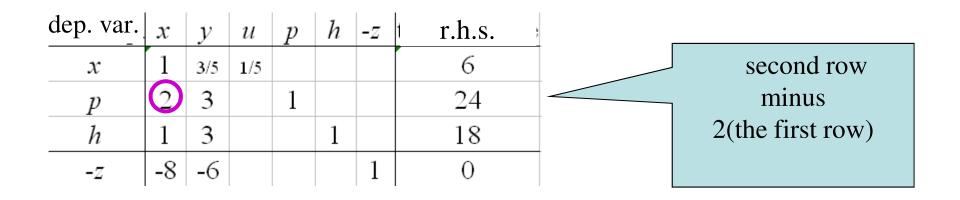
	dep. var.	x	у	и	р	h	- <i>Z</i>	r.h.s.
$x + 3/5y + 1/5u = 6 \longrightarrow$	x	1	3/5	1/5				6
	р	2	3		1			24
	h	1	3			1		18
	-Z	-8	-6				1	0

dep. var.	x	У	и	р	h	- <i>Z</i>	r.h.s.
x	1	3/5	1/5				6
р	2	3		1			24
h	1	3			1		18
-Z	-8	-6				1	0

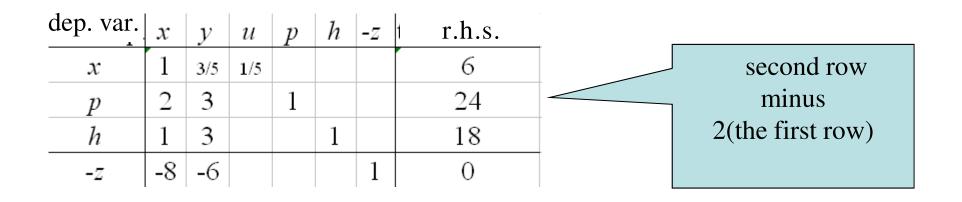
• Recall: Replace *x* in the other equations

$$p = 24 - 2x - 3y$$

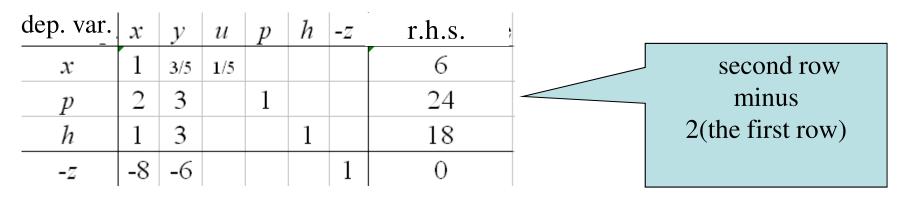
=> $p = 24 - 2(6 - 1/5u - 3/5y) - 3y$
This is equivalent to : $p = 24 - 2(6 - 1/5u - 3/5y) + 2x - 2x - 3y$
 $\Leftrightarrow 2x + 3y + p - 2(x + 3/5y + 1/5u) = 24 - 2(6)$



This is equivalent to : p = 24 - 2(6 - 1/5u - 3/5y) + 2x - 2x - 3y $\Leftrightarrow 2x + 3y + p - 2(x + 3/5y + 1/5u) = 24 - 2(6)$ $\Leftrightarrow 2x + 3y + p = 24$ -2(x + 3/5y + 1/5u) = 60x + 9/5y - 2/5u + p = 12



0x + 9/5y - 2/5u + p = 12



The tableau is modified as follows

	dep. var.	x	У	и	р	h	- z	r.h.s.
	x	1	3/5	1/5				6
0x + 9/5y - 2/5u + p = 12	p		9/5	-2/5	1			12
	h	1	3			1		18
	-Z	-8	-6				1	0

dep. var.	x	y	и	р	h	- <i>z</i>	r.h.s.	
x	1	3/5	1/5				6	
р	2	3		1			24	
h	1	3			1		18	
-Z	-8	-6				1	0	

Doing this for the other rows of the tableau

dep. var.	x	y	и	р	h	- Z	t r.h.s.
x	1	3/5	1/5				6
р		9/5	-2/5	1			12
h		12/5	-1/5		1		12
-Z		-6/5	8/5			1	48

Tableau format of the simplex method analysis of one iteration

• Analyse one iteration of the tableau format of the simplex method

· ·

. .

• The system

•

•

$$\begin{array}{rcl} x_{1}+&&+\overline{a}_{1m+1}x_{m+1}+\ldots+\overline{a}_{1s}x_{s}+\ldots+\overline{a}_{1n}x_{n}=\overline{b}_{1}\\ &&&\\ x_{2}+&&+\overline{a}_{2m+1}x_{m+1}+\ldots+\overline{a}_{2s}x_{s}+\ldots+\overline{a}_{2n}x_{n}=\overline{b}_{2}\end{array}$$

$$x_r + + \bar{a}_{rm+1}x_{m+1} + \dots + \bar{a}_{rs}x_s + \dots + \bar{a}_{rn}x_n = \bar{b}_r$$

$$x_{m} + \overline{a}_{mm+1}x_{m+1} + \dots + \overline{a}_{ms}x_{s} + \dots + \overline{a}_{mn}x_{n} = \overline{b}_{m}$$

$$\overline{c}_{m+1}x_{m+1} + \dots + \overline{c}_{s}x_{s} + \dots + \overline{c}_{n}x_{n} = z - \overline{z}$$

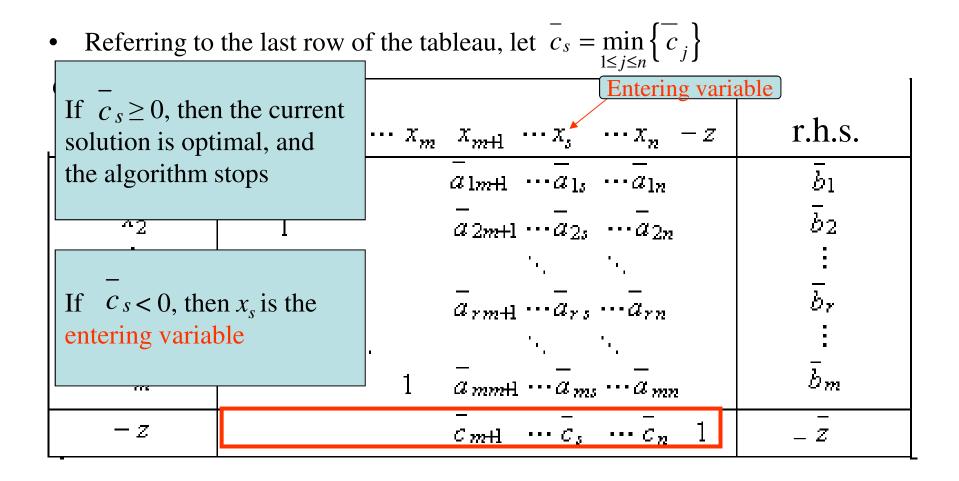
٠

$$\begin{array}{rcl} x_{1}+&&&&+\overline{a}_{1m+1}x_{m+1}+\ldots+\overline{a}_{1s}x_{s}+\ldots+\overline{a}_{1n}x_{n}=\overline{b}_{1}\\ x_{2}+&&&+\overline{a}_{2m+1}x_{m+1}+\ldots+\overline{a}_{2s}x_{s}+\ldots+\overline{a}_{2n}x_{n}=\overline{b}_{2}\\ \cdot&&&\cdot&&\cdot\\ &&&&&\cdot\\ &&&&&&\\ x_{r}+&&+\overline{a}_{rm+1}x_{m+1}+\ldots+\overline{a}_{rs}x_{s}+\ldots+\overline{a}_{rn}x_{n}=\overline{b}_{r}\\ \cdot&&&&&\\ &&&&&\\ x_{m}+\overline{a}_{mm+1}x_{m+1}+\ldots+\overline{a}_{ms}x_{s}+\ldots+\overline{a}_{mn}x_{n}=\overline{b}_{m}\\ &&&&\\ &&&&\\ &&&&\\ &&&&\\ &&&&\\ &&&&\\ &&&&\\ &&&&\\ &&&&\\ &&&&\\ &&&&\\ &&&&\\ &&&\\ &&&&\\ &&&\\$$

can be written in the following tableau

dep. var.		
	$x_1 \ x_2 \cdots x_r \cdots x_m \ x_{m+1} \ \cdots x_s \ \cdots x_n \ -z$	r.h.s.
<i>x</i> ₁	1 $\overline{a}_{1m+1} \cdots \overline{a}_{1s} \cdots \overline{a}_{1n}$	\overline{b}_1
x ₂	1 $\overline{a}_{2m+1}\cdots\overline{a}_{2s}\cdots\overline{a}_{2n}$	\overline{b}_2
:	··. ··. ··.	:
<i>x</i> ,	1 $\overline{a}_{rm+1} \cdots \overline{a}_{rs} \cdots \overline{a}_{rn}$	b,
:	··. ··. ··.	:
<i>x</i> _m	$1 \overline{a}_{mmH} \cdots \overline{a}_{ms} \cdots \overline{a}_{mn}$	b_m
- <i>z</i>	$\bar{c}_{m+1} \cdots \bar{c}_s \cdots \bar{c}_n = 1$	$-\overline{z}$

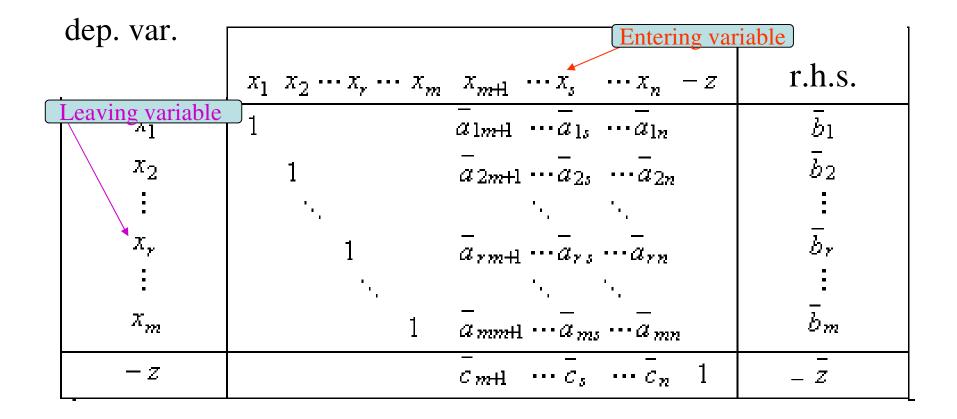
Step1: Select the entering variable

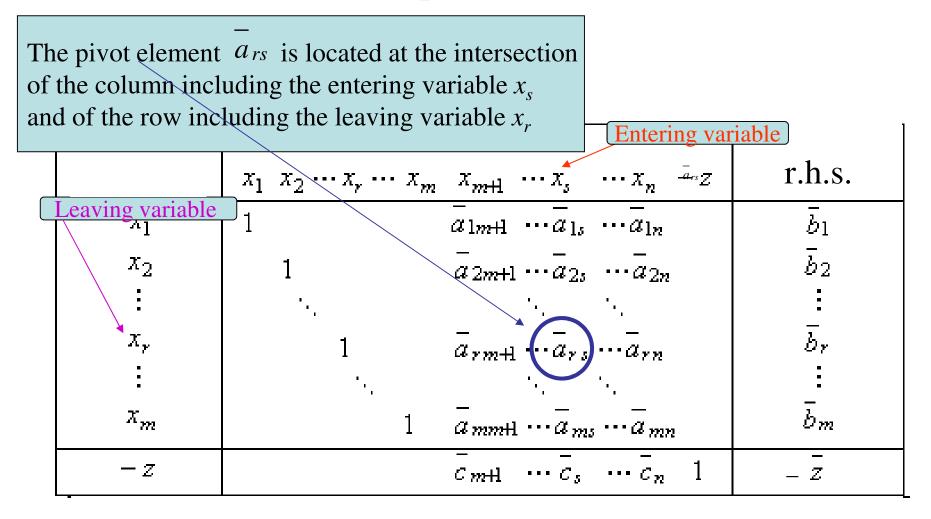


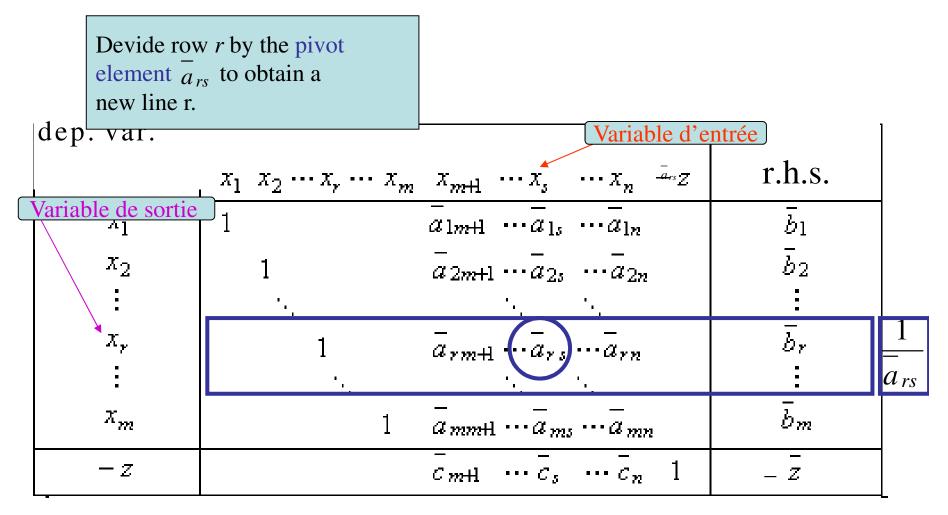
Step 2: Select the leaving variable

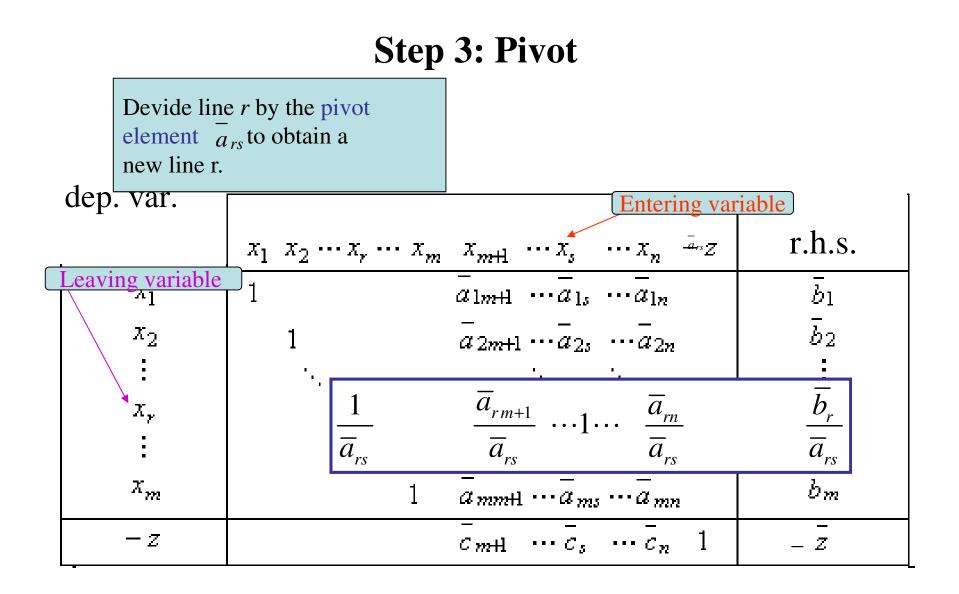
If $\overline{a}_{is} \leq 0 \forall \ 1 \leq i \leq m$, the problem is not	$\cdot x_r \cdots x_m x_m$		Entering var	riable r.h.s.
	_			
and the alg. stops	a_{1n}	$\mathbf{H} \cdot \mathbf{a}_{1s}$		b_1
	a 2,	$n+1 \cdot \cdot a_{2s}$	$\cdots \overline{a}_{2n}$	\overline{b}_2
If $\exists i \text{ such that } \overline{a}_{is} > 0$,	_	· -	·., —	$\frac{1}{\overline{b}_r}$
then the sol. remains feasi	ble \Leftrightarrow a_{γ}	$n+1 \cdots a_{rs}$	$\cdots a_{rn}$:
$\forall i \text{ such that } \overline{a_{is}} > 0$		-	'. 	\overline{b}_m
	—	mti • a mi		
$x_i = \overline{b}_i - \overline{a}_{is} x_s \ge 0 \Leftrightarrow $	$x_s \leq \frac{\sigma_i}{\sigma_i}$	The entering	ng variable.	x_s takes the value
	a_{is}	$x = \frac{\overline{l}}{2}$	$\frac{\overline{b}_r}{a_{rs}} = \min_{1 \le i \le m} \left\{ \frac{\overline{b}}{a} \right\}$	$\left\{\begin{array}{c} -\\ -\\ -\\ -\\ -\\ -\\ -\\ -\\ -\\ -\\ -\\ -\\ -\\ $
		\mathcal{L}_{s}	$l_{rs} = \frac{1 \le i \le m}{1 \le i \le m} \left[\frac{1}{a} \right]$	is

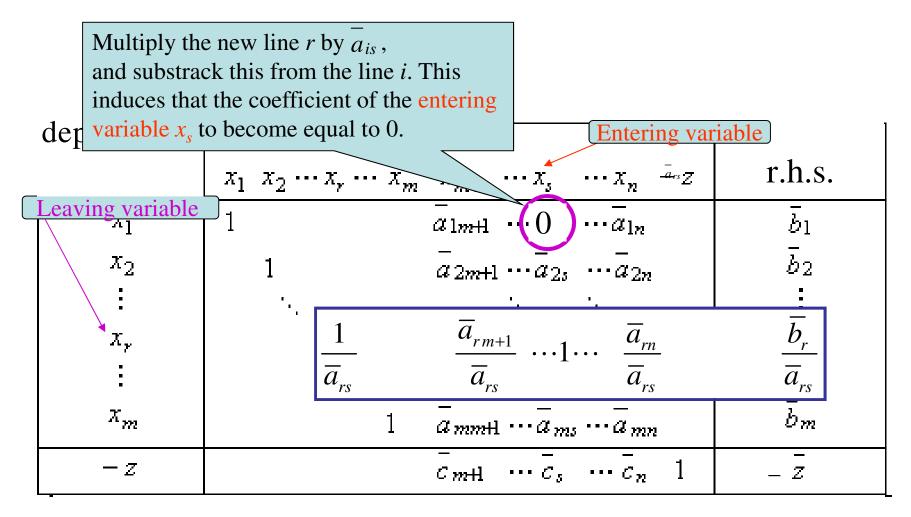
Step 2: Select the leaving variable

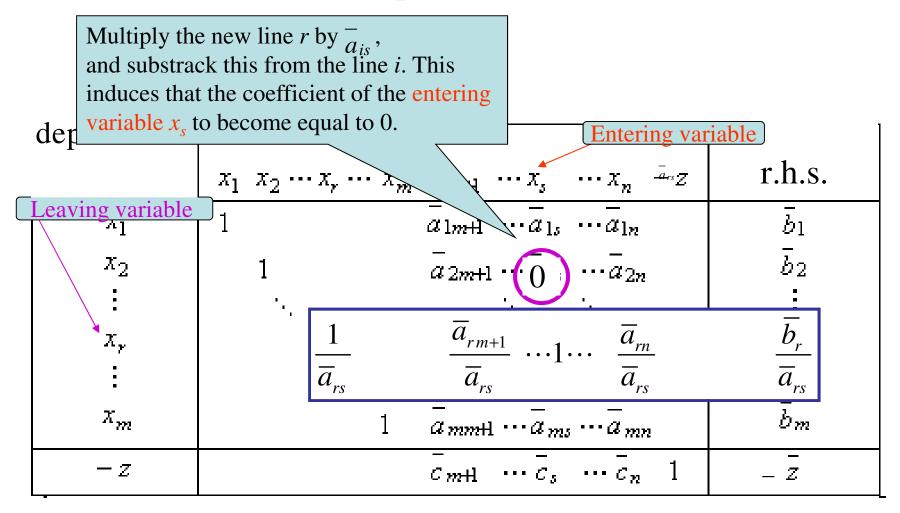


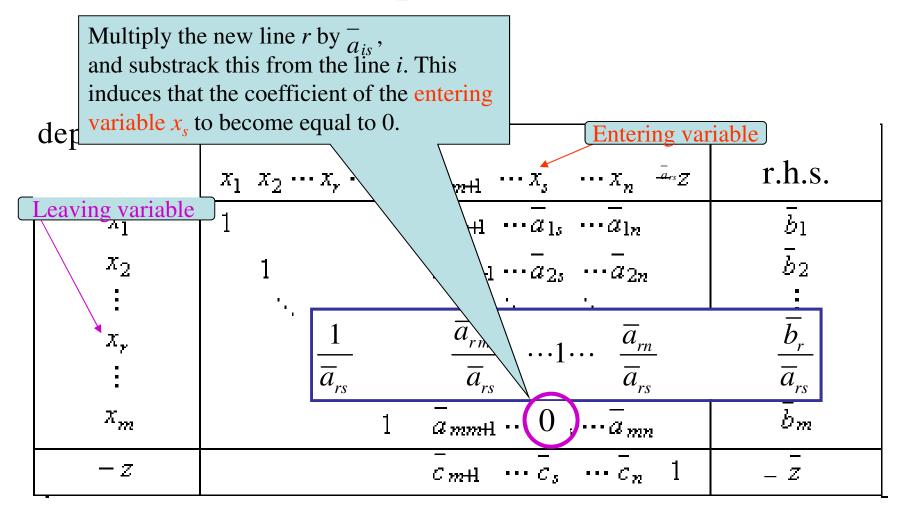


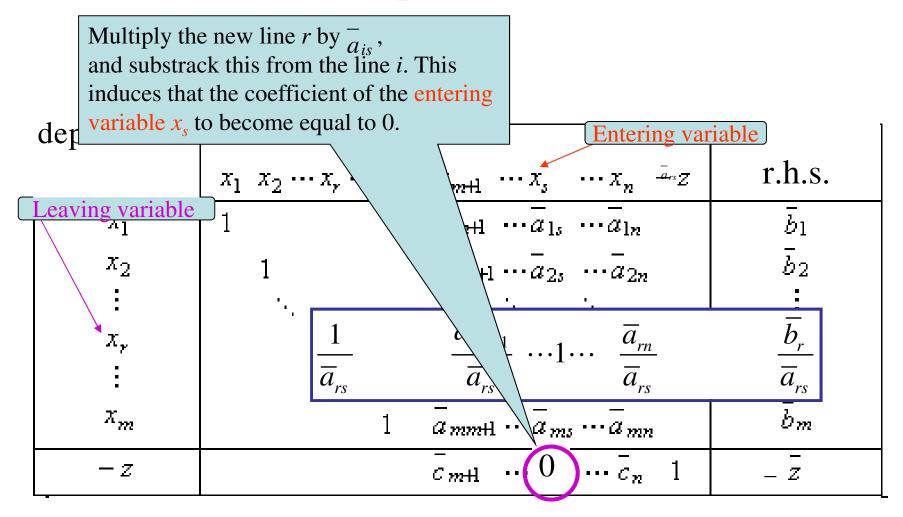












New tableau for the next iteration

dep. var.		-
	$x_1 x_2 \cdots x_r \cdots x_m x_{m+1} \cdots x_s \cdots x_n - z$	r.h.s.
<i>x</i> ₁	1 $\cdots \overline{\widetilde{a}}_{1r} \cdots \overline{\widetilde{a}}_{1m+1} \cdots 0 \cdots \overline{\widetilde{a}}_{1n}$	$\overline{\widetilde{b}}_1$
x ₂	$1 \cdots \overline{\widetilde{a}}_{2r} \cdots \qquad \overline{\widetilde{a}}_{2m+1} \cdots 0 \cdots \overline{\widetilde{a}}_{2n}$	$\overline{\tilde{b}}_2$
		:
x _s	$\cdots \overline{\widetilde{a}}_{rr} \cdots \overline{\widetilde{a}}_{rm+1} \cdots 1 \cdots \overline{\widetilde{a}}_{rn}$	$\overline{\widetilde{b}}_r$
	'·. '·. '·. 	-
x _m	$\cdots \overline{\widetilde{a}}_{mr} \cdots 1 \ \overline{\widetilde{a}}_{mm+1} \cdots 0 \cdots \overline{\widetilde{a}}_{mn}$	$\overline{\widetilde{b}}_m$
- <i>z</i>	$\cdots \overline{\widetilde{c}}_{r} \cdots \overline{\widetilde{c}}_{m+1} \cdots 0 \cdots \overline{\widetilde{c}}_{n} = 1$	$-\overline{\widetilde{z}}$

• The linear programming problem in standard form

min
$$z = c_1 x_1 + c_2 x_2 + ... + c_n x_n$$

Subject to $a_{11} x_1 + a_{12} x_2 + ... + a_{1n} x_n = b_1$
 $a_{21} x_1 + a_{22} x_2 + ... + a_{2n} x_n = b_2$
 $...$
 $a_{m1} x_1 + a_{m2} x_2 + ... + a_{mn} x_n = b_m$
 $x_1, x_2, ..., x_n \ge 0$

Restaurant owner problem min $z = [-8, -6, 0, 0, 0] \begin{bmatrix} x \\ y \end{bmatrix}$ y и р h s.t. ${\mathcal X}$ $\begin{pmatrix} 5 & 3 & 1 & 0 & 0 \\ 2 & 3 & 0 & 1 & 0 \\ 1 & 3 & 0 & 0 & 1 \end{pmatrix} \begin{vmatrix} y \\ u \\ p \end{vmatrix} = \begin{bmatrix} 30 \\ 24 \\ 18 \end{bmatrix}$ $x, y, u, p, h \ge 0$

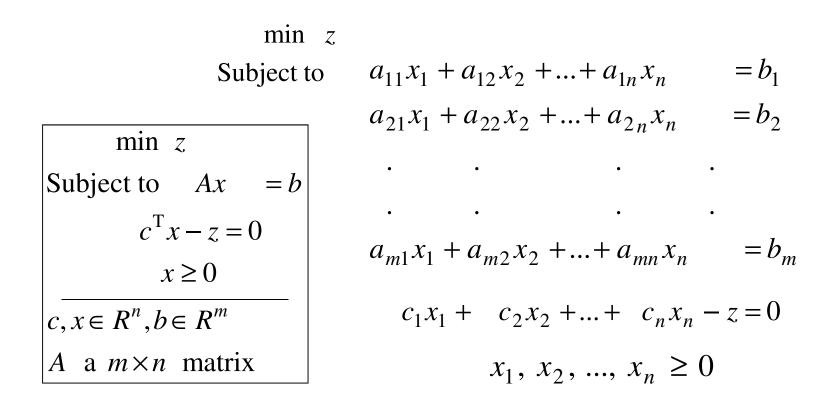
min
$$z = -8x - 6y$$

Suject to $5x + 3y + u = 30$
 $2x + 3y + p = 24$
 $1x + 3y + h = 18$
 $x, y, u, p, h \ge 0$

• The linear programming problem in standard form

$$\begin{array}{rcl}
& \min & z = c_1 x_1 + c_2 x_2 + \ldots + c_n x_n \\
& \text{Subject to} & a_{11} x_1 + a_{12} x_2 + \ldots + a_{1n} x_n = b_1 \\
\\
& \prod_{\substack{n \neq n \\ n \neq n}} z = c^T x \\
& \text{subject to} & Ax = b \\
& x \ge 0 \\
\hline
& c, x \in \mathbb{R}^n, b \in \mathbb{R}^m \\
& A \ am \times n \ matrix
\end{array}$$

$$\begin{array}{rcl}
& a_{21} x_1 + a_{22} x_2 + \ldots + a_{2n} x_n = b_2 \\
& \ddots & \ddots & \ddots \\
& \ddots & \ddots & \ddots \\
& a_{m1} x_1 + a_{m2} x_2 + \ldots + a_{mn} x_n = b_m \\
& x_1, x_2, \dots, x_n \ge 0 \\
\hline
& A = \begin{pmatrix} a_{11} & \ldots & a_{1n} \\ \vdots & \ddots & \vdots \\ a_{m1} & \cdots & a_{mn} \end{pmatrix} \quad b = \begin{bmatrix} b_1 \\ \vdots \\ b_m \end{bmatrix} \\
& c^T = [c_1, \dots, c_n] \\
\end{array}$$



• Consider the matrix formulation of the linear programming problem

```
min z
Subject to Ax = b
c^{T}x - z = 0
x \ge 0
```

- Assume that $m \le n$ and that the matrix A is of full rank (i.e., rank(A) = m, or that the rows of A are linearly independent)
- A sub matrix *B* of *A* is a basis of *A* if it is a *mxm* matrix and non singular (i.e, *B*⁻¹ exists)

Restaurant owner problem:

Examples of basis:

$$u \ p \ h \qquad x \ p \ h \qquad x \ p \ y$$
$$B_{1} = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix} \qquad B_{2} = \begin{pmatrix} 5 & 0 & 0 \\ 2 & 1 & 0 \\ 1 & 0 & 1 \end{pmatrix} \qquad B_{3} = \begin{pmatrix} 5 & 0 & 3 \\ 2 & 1 & 3 \\ 1 & 0 & 3 \end{pmatrix}$$

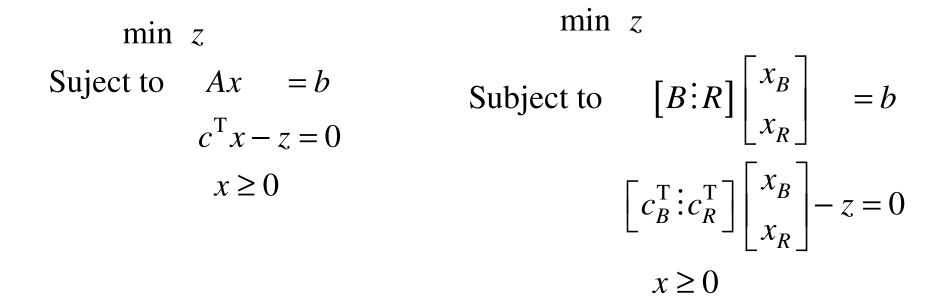
- A sub matrix *B* of *A* is a basis of *A* if it is a *mxm* matrix and non singular (i.e, *B*⁻¹ exists)
- To ease the presentation, assume that the basis *B* includes the first *m* columns of *A*, and then

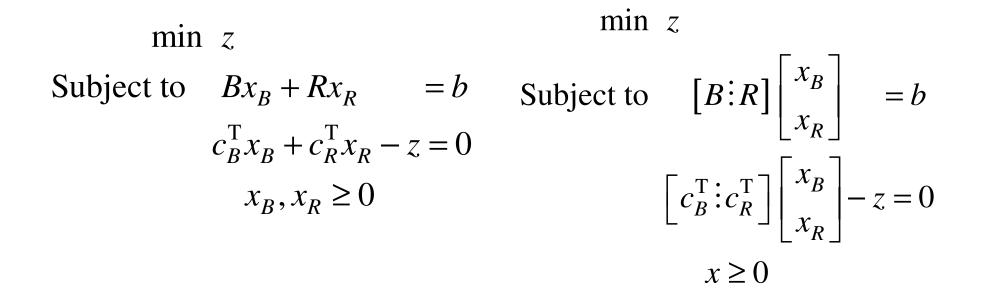
$$A = \begin{bmatrix} B \vdots R \end{bmatrix}$$

Denote also

$$x = \begin{bmatrix} x_B \\ x_R \end{bmatrix} \qquad \qquad c = \begin{bmatrix} c_B \\ c_R \end{bmatrix}$$

• The original problem can be written as



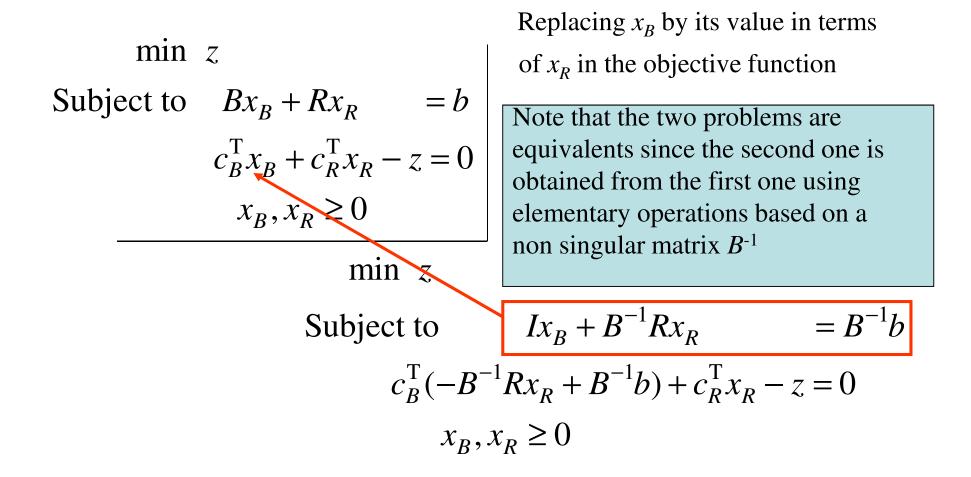


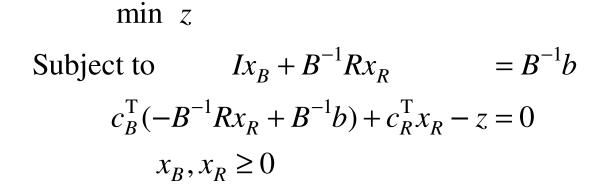
• Specify x_B in terms of x_R using the constraints of the problem

$$Bx_B + Rx_R = b$$
$$B^{-1}(Bx_B + Rx_R) = B^{-1}b$$
$$B^{-1}Bx_B + B^{-1}Rx_R = B^{-1}b$$
$$Ix_B + B^{-1}Rx_R = B^{-1}b$$

• Then

$$Ix_B = -B^{-1}Rx_R + B^{-1}b$$

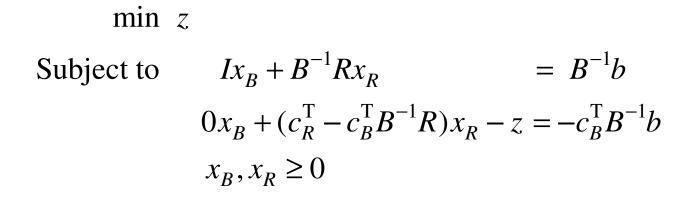




Combining the coefficients of x_R

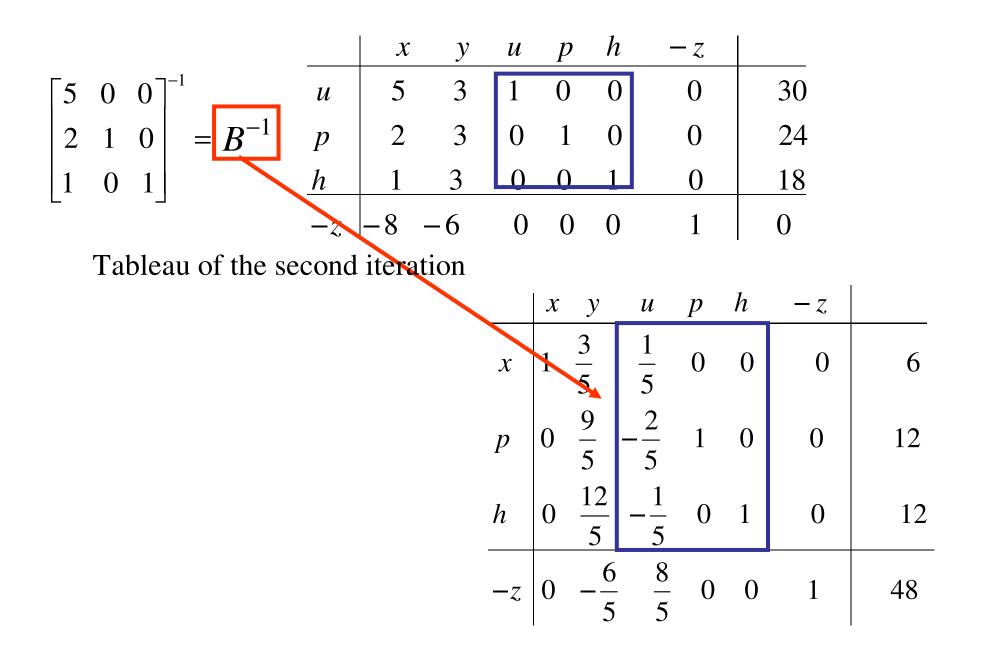
min z

Subject to $Ix_B + B^{-1}Rx_R = B^{-1}b$ $0x_B + (c_R^T - c_B^T B^{-1}R)x_R - z = -c_B^T B^{-1}b$ $x_B, x_R \ge 0$



The problem can be specified in the following tableau

dep. var.	x_B^T	x_R^T	- <i>z</i>	r.h.s.
X _B	Ι	$B^{-1}R$	0	$B^{-1}b$
<i>Z</i>	0	$c_R^T - c_B^T B^{-1} R$! 1	$-c_B^T B^{-1}b$



basic var.	x_B^T	x_R^T	- z	r.h.s.
x_B	Ι	$B^{-1}R$	0	$B^{-1}b$
<i>Z</i>	0	$c_R^T - c_B^T B^{-1} R$	1	$-c_B^T B^{-1} b$

The variables x_B (denoted as the dependent variables) associated with the columns of the basis *B*, are now denoted basic variables

The variables x_R (denoted independent variables) are now denoted non basic variables

basic var.	x_B^I	x_R^T	- <i>z</i>	r.h.s.
x_B	Ι	$B^{-1}R$	0	$B^{-1}b$
<i>z</i>	0	$c_R^T - c_B^T B^{-1} F$	2 1	$-c_B^T B^{-1} b$

To obtain the basic solution associated with the basis B let $x_R = 0$ it follows that $x_B = B^{-1}b$ The basic solution is feasible if $x_B \ge 0$.

basic var.	x_B^T	x_R^T	- <i>z</i>	r.h.s.
x_B	Ι	$B^{-1}R$	0	$B^{-1}b$
<i>Z</i>	0	$c_R^T - c_B^T B^{-1} i$	२ 1	$-c_B^T B^{-1} b$

Note that this tableau is identical to the one used to illustrate an iteration of the simplex algorithm.

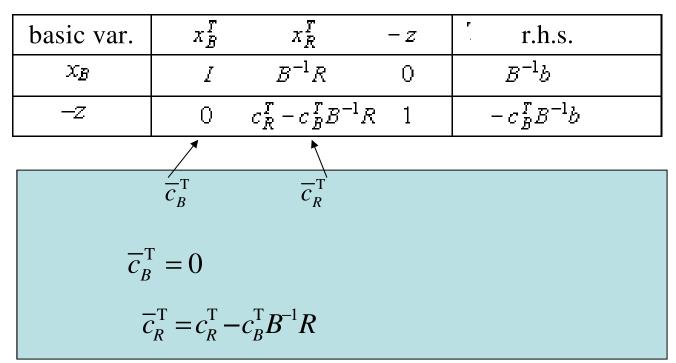
basic var.	x_B^I	x_R^T	- <i>z</i>	r.h.s.
χ_B	Ι	$B^{-1}R$	0	$B^{-1}b$
Z	0	$c_R^T - c_B^T B^{-1} R$	1	$-c_B^T B^{-1} b$

dep. var.		r.h.s.
	$x_1 x_2 \cdots x_r \cdots x_m x_{m+1} \cdots x_s \cdots x_n - z$	
$ \begin{array}{c} x_1 \\ x_2 \\ \vdots \\ x_r \\ \vdots \\ x_m \end{array} $	$\begin{bmatrix} 1 & & \overline{a_{1m+1}} & \cdots & \overline{a_{1s}} & \cdots & \overline{a_{1n}} \\ 1 & & \overline{a_{2m+1}} & \cdots & \overline{a_{2s}} & \cdots & \overline{a_{2n}} \\ & & \ddots & & \ddots & & \ddots \\ & & & & & & \ddots & \ddots$	
- <i>z</i>	$\overline{c_{m+1}} \cdots \overline{c_s} \cdots \overline{c_n} $	- Z

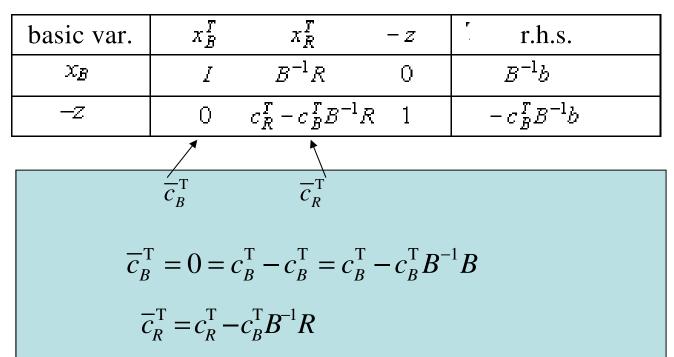
basic var.	x_B^I	x_R^T	- <i>z</i>	r.h.s.
x_B	Ι	$B^{-1}R$	0	$B^{-1}b$
<i>z</i>	0	$c_R^T - c_B^T B^{-1} R$	2 1	$-c_B^T B^{-1} b$

Since all tableaus correspond to basis of *A* that are specified by the columns associated with the basic variables (dependent variables), it follows that in the simplex algorithm, we move from one basic feasible solution to a new one where the objective function decreases.

• Consider the last row of the simplex tableau corresponding to the basis *B* associated to the relative costs of the variables:



• Consider the last row of the simplex tableau corresponding to the basis *B* associated to the relative costs of the variables:



$$\overline{c}^{\mathrm{T}} = \left[\overline{c}_{B}^{\mathrm{T}}, \overline{c}_{R}^{\mathrm{T}}\right] = \left[c_{B}^{\mathrm{T}}, c_{R}^{\mathrm{T}}\right] - c_{B}^{T}B^{-1}\left[B \colon R\right] = c^{\mathrm{T}} - c_{B}^{T}B^{-1}A$$

$$\overline{c}^{\mathrm{T}} = c^{\mathrm{T}} - c_{B}^{\mathrm{T}} B^{-1} A$$

Denote the vector $\pi \in \mathbb{R}^m$ specified by

$$\pi^{\mathrm{T}} = c_{B}^{\mathrm{T}} B^{-1}$$
Then

$$\overline{c}^{\mathrm{T}} = c^{\mathrm{T}} - \pi^{\mathrm{T}} A$$

or

$$\begin{bmatrix} \overline{c}_1, \dots, \overline{c}_n \end{bmatrix} = \begin{bmatrix} c_1, \dots, c_n \end{bmatrix} - \pi^{\mathrm{T}} \begin{bmatrix} a_{\bullet 1}, \dots, a_{\bullet n} \end{bmatrix}$$
$$\overline{c}_j = c_j - \pi^{\mathrm{T}} a_{\bullet j}$$

where $a_{\bullet j}$ denotes the *jth* column of the contraint matrix A

 π is the simplex multipliers vector associated with the basis *B*.

The vector π has one element associated with each row (constraint) of the tableau.

	X	<i>y</i>	U	p	h	<u> </u>		$\boldsymbol{\pi}^{\mathrm{T}} = \boldsymbol{c}_{B}^{\mathrm{T}}\boldsymbol{B}^{-1}$
X	1	$\frac{3}{5}$	$\frac{1}{5}$	0	0	0	6	□
p	0	$\frac{9}{5}$	$-\frac{2}{5}$	1	0	0	12	$\pi^{\mathrm{T}} = \begin{bmatrix} -8 & 0 & 0 \end{bmatrix} \qquad \begin{bmatrix} \frac{1}{5} \\ -\frac{2}{5} \\ -\frac{1}{5} \end{bmatrix}$
h	0	$\frac{12}{5}$	$-\frac{1}{5}$	0	1	0	12	$\begin{bmatrix} -\frac{1}{5} \end{bmatrix}$
-z	0	$-\frac{6}{5}$	$\frac{8}{5}$	0	0	1	48	$\overline{c}^{\mathrm{T}} = c^{\mathrm{T}} - \pi^{\mathrm{T}} A$
								$\pi^{\mathrm{T}}A = \left[-\frac{8}{5} \ 0 \ 0\right] \begin{bmatrix} 5\\2\\1 \end{bmatrix}$
		У	U	Į)	h - z		
U	5	3	1	()	0 0	30	$=\left[-8 - \frac{24}{5} - \frac{8}{5} \right] 0$
р	2	3	0		1	0 0	24	$\overline{c}^{\mathrm{T}} = c^{\overline{\mathrm{T}}} - \pi^{\mathrm{T}} A$
h	1	3	0	(С	1 0	18	= [-8 -6 0 0 0] -
<i>-z</i>	-8 ↑	-6	0) ()	0 1 ↑	0	$= \left[0 \ -\frac{6}{5} \ \frac{8}{5} \ 0 \ 0 \right]$

 $\begin{bmatrix} 0 & 0 \\ 1 & 0 \\ 0 & 1 \end{bmatrix} = \begin{bmatrix} -\frac{8}{5} & 0 & 0 \end{bmatrix}$

 $\begin{bmatrix} 3 & 1 & 0 & 0 \\ 3 & 0 & 1 & 0 \\ 3 & 0 & 0 & 1 \end{bmatrix}$ 0 0 $\left[-8 \ -\frac{24}{5} \ -\frac{8}{5} \ 0 \ 0\right]$

hp -zX У U $\pi^{\mathrm{T}} = c_{B}^{\mathrm{T}} B^{-1}$ $-\frac{1}{4}$ $-\frac{3}{4}$ $\frac{5}{12}$ $\frac{1}{4}$ $\pi^{\mathrm{T}} = \begin{bmatrix} -8 & 0 & -6 \end{bmatrix} \begin{bmatrix} \frac{1}{4} & 0 & -\frac{1}{4} \\ -\frac{1}{4} & 1 & -\frac{3}{4} \\ -\frac{1}{12} & 0 & \frac{5}{12} \end{bmatrix} = \begin{bmatrix} -\frac{3}{2} & 0 & -\frac{1}{2} \end{bmatrix}$ 0 x 0 0 3 $-\frac{1}{4}$ 1 3 0 0 p 0 $-\frac{1}{12}$ 0 3 1 y 0 0 $\frac{3}{2}$ $\frac{1}{2}$ 0 0 $-z \mid 0$ 1 54 $\overline{c}^{\mathrm{T}} = c^{\mathrm{T}} - \pi^{\mathrm{T}} A$ $\pi^{\mathrm{T}} A = \begin{bmatrix} -\frac{3}{2} & 0 & -\frac{1}{2} \end{bmatrix} \begin{bmatrix} 5 & 3 & 1 & 0 & 0 \\ 2 & 3 & 0 & 1 & 0 \\ 1 & 3 & 0 & 0 & 1 \end{bmatrix}$ $= \begin{bmatrix} -8 & -6 & -\frac{3}{2} & 0 & -\frac{1}{2} \end{bmatrix}$ -zр Х y \mathcal{U} 3 5 0 1 0 0 U 2 3 0 0 0 1 p $18 \quad \overline{c}^{\mathrm{T}} = c^{\mathrm{T}} - \pi^{\mathrm{T}} A$ h 1 3 0 0 1 0 $\frac{18}{0} = \begin{bmatrix} -8 & -6 & 0 & 0 \end{bmatrix} - \begin{bmatrix} -8 & -6 & -\frac{3}{2} & 0 & -\frac{1}{2} \end{bmatrix}$ -z-8 -6 0 1 0 0 1 $= \begin{bmatrix} 0 & 0 & \frac{3}{2} & 0 & \frac{1}{2} \end{bmatrix}$

 $\bar{c}_j = c_j - \pi^{\mathrm{T}} a_{\bullet j}$

The simplex multipliers vector π alows to compute the values of the relative costs \overline{c}_i using the original data of the problem.

$$\overline{c}^{\mathrm{T}} = c^{\mathrm{T}} - \pi^{\mathrm{T}} A$$

The component π_i (i = 1,...,m) of the simplex multiplier vector can be seen as a weight associated with the row *i* of the constraint matrix *A*, and then substracting the linear combination of the rows of the matrix *A* using these weights $(\pi^T A)$, from the cost vector (c^T) of the variables allows to obtain the relative costs of the variables, and to reduce to 0 those of the basic variables.

h р -zX V U $\pi^{\mathrm{T}} = c_{B}^{\mathrm{T}} B^{-1}$ $-\frac{1}{4}$ $-\frac{3}{4}$ $\frac{5}{12}$ $\frac{1}{4}$ $\pi^{\mathrm{T}} = \begin{bmatrix} -8 & 0 & -6 \end{bmatrix} \begin{bmatrix} \frac{1}{4} & 0 & -\frac{1}{4} \\ -\frac{1}{4} & 1 & -\frac{3}{4} \\ -\frac{1}{12} & 0 & \frac{5}{12} \end{bmatrix} = \begin{bmatrix} -\frac{3}{2} & 0 & -\frac{1}{2} \end{bmatrix}$ 0 x 0 0 3 $-\frac{1}{4}$ 1 0 3 0 0 p $-\frac{1}{12}$ 0 3 y 1 0 0 $\frac{3}{2}$ $\frac{1}{2}$ 0 0 $-z \mid 0$ 1 54 $\overline{c}^{\mathrm{T}} = c^{\mathrm{T}} - \pi^{\mathrm{T}} A$ $\pi^{\mathrm{T}}A = \begin{bmatrix} -\frac{3}{2} & 0 & -\frac{1}{2} \end{bmatrix} \begin{bmatrix} 5 & 3 & 1 & 0 & 0 \\ 2 & 3 & 0 & 1 & 0 \\ 1 & 3 & 0 & 0 & 1 \end{bmatrix}$ $30 = \begin{bmatrix} -8 & -6 & -\frac{3}{2} & 0 & -\frac{1}{2} \end{bmatrix}$ -zр h \mathcal{U} Х 3 5 -3/2 1 0 0 0 2 3 0 0 1 0 0 $\frac{18}{0} = \begin{bmatrix} -8 & -6 & 0 & 0 \end{bmatrix} - \begin{bmatrix} -8 & -6 & -\frac{3}{2} & 0 & -\frac{1}{2} \end{bmatrix}$ $\frac{-1/2}{-z}$ 1 3 0 0 1 0 -8 -6 0 0 0 1 $= \begin{bmatrix} 0 & 0 & \frac{3}{2} & 0 & \frac{1}{2} \end{bmatrix}$

- The simplex multipliers associated with an optimal basic solution allows to analyse the sensitivity of the optimal value when modfying the right hand terms.
- Consider a problem and its version when modfying the right hand terms

$$\begin{array}{c|cccc} \min z & \min \tilde{z} \\ \text{Subject to} & Ax &= b \\ c^{\mathrm{T}}x - z &= 0 \\ x \geq 0 \end{array} & \begin{array}{c|cccccc} \min \tilde{z} \\ \text{Subject to} & A\tilde{x} &= b + \Delta b \\ \text{Subject to} & A\tilde{x} &= b + \Delta b \\ c^{\mathrm{T}}\tilde{x} - \tilde{z} &= 0 \\ \tilde{x} \geq 0 \end{array}$$

1

$$\begin{array}{c|c} \min z \\ \text{Subject to} \quad Ax = b \\ c^{\mathrm{T}}x - z = 0 \\ x \ge 0 \end{array} \qquad \begin{array}{c|c} \min \tilde{z} \\ \text{Subject to} \quad A\tilde{x} = b + \Delta b \\ \text{Subject to} \quad A\tilde{x} = b + \Delta b \\ c^{\mathrm{T}}\tilde{x} - \tilde{z} = 0 \\ \tilde{x} \ge 0 \end{array}$$

Let B^* be an optimal basis, and the corresponding basic solution

$$x_{R}^{*} = 0$$

$$x_{B^{*}}^{*} = B^{*-1}b = \overline{b} \ge 0$$

having the optimal value

$$z^{*} = c_{B^{*}}^{T} x_{B^{*}}^{*} + c_{R}^{T} x_{R}^{*} = c_{B^{*}}^{T} B^{*-1} b = c_{B^{*}}^{T} \overline{b}$$

$$\begin{array}{c|c} \min z \\ \text{Subject to} \quad Ax = b \\ c^{\mathrm{T}}x - z = 0 \\ x \ge 0 \end{array} \qquad \begin{array}{c|c} \min \tilde{z} \\ \text{Subject to} \quad A\tilde{x} = b + \Delta b \\ \text{Subject to} \quad A\tilde{x} = b + \Delta b \\ c^{\mathrm{T}}\tilde{x} - \tilde{z} = 0 \\ \tilde{x} \ge 0 \end{array}$$

T

Select a value of the vector Δb to be such that

$$B^{*-1}(b + \Delta b) = B^{*-1}b + B^{*-1}\Delta b \ge 0$$

Hence B^* remains a feasible basis for the new problem, and the basic feasible solution is as follows

$$\widetilde{x}_{R}^{*} = 0$$

$$\widetilde{x}_{B^{*}}^{*} = B^{*-1}(b + \Delta b) \ge 0$$

Hence B^* remains a feasible basis for the new problem, and the basic feasible solution is as follows $\overline{c}^{*T} = c^{T}$.

$$\widetilde{x}_{R}^{*} = 0$$

$$\widetilde{x}_{B^{*}}^{*} = B^{*-1}(b + \Delta b) \ge 0$$

$$\overline{c}^{*\mathrm{T}} = c^{\mathrm{T}} - \pi^{*\mathrm{T}}A$$

$$\boldsymbol{\pi}^{*\mathrm{T}} = \boldsymbol{c}_{\boldsymbol{B}^*}^{\mathrm{T}} \boldsymbol{B}^{*-1}$$

Furthermore, since neither the costs c_j nor the matrix A have changed, then the simplex multipliers vector π^* remains unchanged. Then the relative costs \overline{c}_j remains also unchanged and non negative for the new problems.

Then B^* remains an optimal basis for the new problem.

An optimal solution of the new problem is then

$$\tilde{x}_{R}^{*} = 0$$

 $\tilde{x}_{B^{*}}^{*} = B^{*-1}(b + \Delta b) \ge 0$

The optimal value of the new problem is as follows: $\tilde{z}^* = c_{B^*}^T \tilde{x}_{B^*}^* + c_R^T \tilde{x}_R^*$

$$= c_{B^{*}}^{T} B^{*-1} (b + \Delta b)$$

= $c_{B^{*}}^{T} B^{*-1} b + c_{B^{*}}^{T} B^{*-1} \Delta b$
= $z^{*} + \pi^{*T} \Delta b$

$$= z^* + \sum_{i=1}^m \pi_i^* \Delta b_i$$

$$\pi^{*\mathrm{T}} = c_{B^*}^{\mathrm{T}} B^{*-1}$$

$$z^* = c_{B^*}^{\mathrm{T}} B^{*-1} b$$

An optimal solution of the new problem is then

$$\tilde{z}^{*} = c_{B^{*}}^{T} \tilde{x}_{B^{*}}^{*} + c_{R}^{T} \tilde{x}_{R}^{*}$$

$$= c_{B^{*}}^{T} B^{*-1} (b + \Delta b)$$

$$= c_{B^{*}}^{T} B^{*-1} b + c_{B^{*}}^{T} B^{*-1} \Delta b$$

$$= z^{*} + \pi^{*T} \Delta b$$

$$= z^{*} + \sum_{i=1}^{m} \pi_{i}^{*} \Delta b_{i}$$

Then π_i^* is the marginal rate of variation of the optimal value of the objective function when the right hand term b_i of the i^{th} constraint is modified in such a way that the optimal basis remains feasible for the new problem.

	X	У	и	р	h	<i>-z</i>		
x	1	0	$\frac{1}{4}$	0	$-\frac{1}{4}$	0	3	
p	0	0	$-\frac{1}{4}$	1	$-\frac{3}{4}$	0	3	$\pi^{\mathrm{T}} = [$
у	0	1	$-\frac{1}{12}$	0	$\frac{5}{12}$	0	3	
- <i>z</i> .	0	0	$\frac{3}{2}$	0	$\frac{1}{2}$	1	54	

$$\pi^{\mathrm{T}} = c_{B}^{\mathrm{T}} B^{-1}$$

$$\pi^{\mathrm{T}} = \begin{bmatrix} -8 & 0 & -6 \end{bmatrix} \begin{bmatrix} \frac{1}{4} & 0 & -\frac{1}{4} \\ -\frac{1}{4} & 1 & -\frac{3}{4} \\ -\frac{1}{12} & 0 & \frac{5}{12} \end{bmatrix} = \begin{bmatrix} -\frac{3}{2} & 0 & -\frac{1}{2} \end{bmatrix}$$

$$\pi^{*T} = c_{B^*}^{T} B^{*-1}$$

$$\pi^{*T} = \begin{bmatrix} -8 & 0 & -6 \end{bmatrix} \begin{bmatrix} \frac{1}{4} & 0 & -\frac{1}{4} \\ -\frac{1}{4} & 1 & -\frac{3}{4} \\ -\frac{1}{12} & 0 & \frac{5}{12} \end{bmatrix} = \begin{bmatrix} -\frac{3}{2} & 0 & -\frac{1}{2} \end{bmatrix}$$

$$\tilde{z}^{*} = z^{*} + \pi^{*T} \Delta b$$

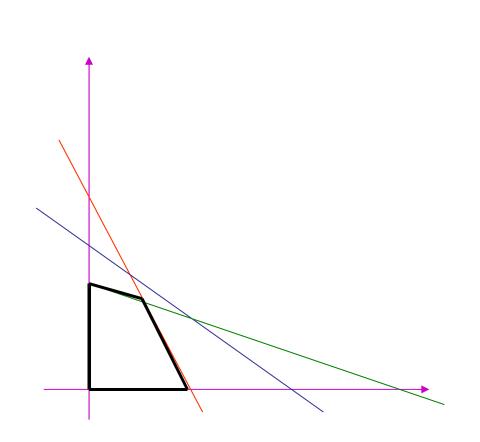
= $-54 + \left[-\frac{3}{2} \ 0 \ -\frac{1}{2} \right] \begin{bmatrix} \Delta b_{1} \\ \Delta b_{2} \\ \Delta b_{3} \end{bmatrix} = -54 - \frac{3}{2} \Delta b_{1} + 0 \Delta b_{2} - \frac{1}{2} \Delta b_{3}$

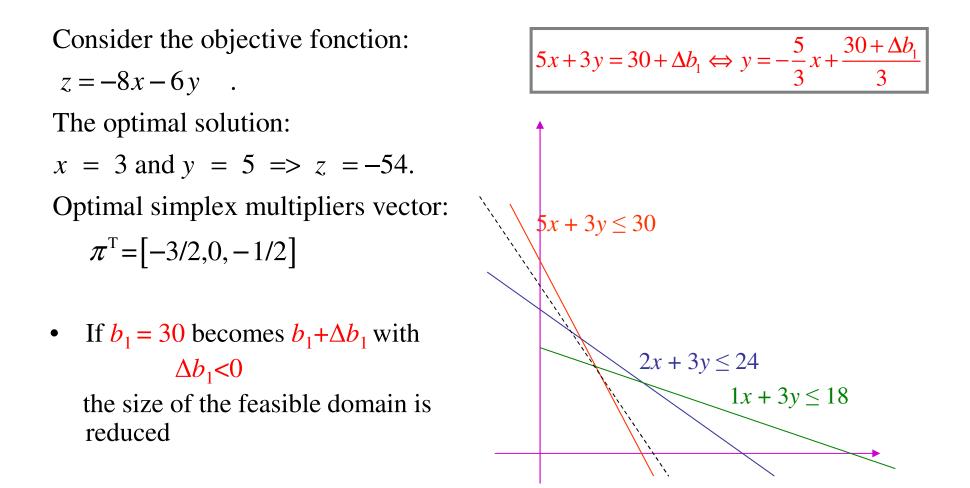
Let
$$\Delta b_2 = \Delta b_3 = 0$$
.
 $\Delta b_1 < 0 \Longrightarrow -\frac{3}{2} \Delta b_1 > 0 \Longrightarrow \tilde{z}^* > z^*$

Feasible domain

• The feasible domain for the system

 $5x + 3y \le 30$ $2x + 3y \le 24$ $1x + 3y \le 18$ $x, y \ge 0$



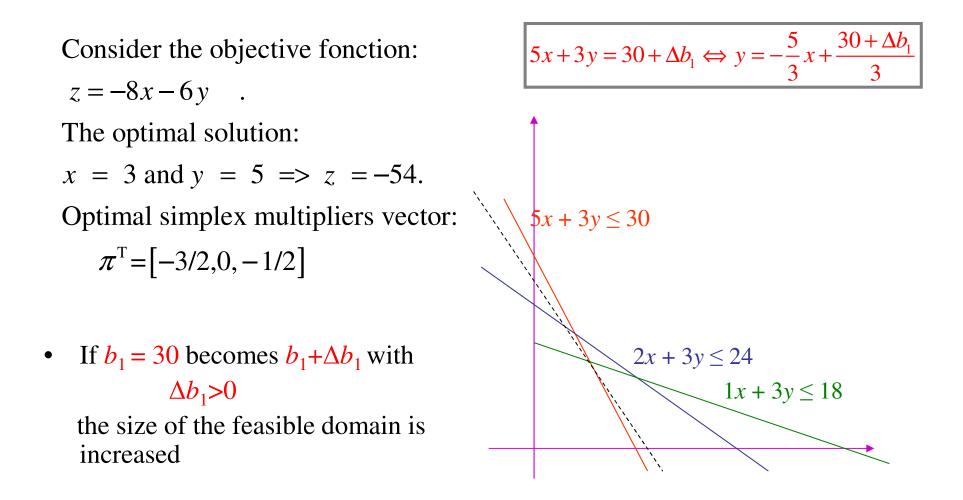


$$\pi^{*T} = c_{B^*}^{T} B^{*-1}$$

$$\pi^{*T} = \begin{bmatrix} -8 & 0 & -6 \end{bmatrix} \begin{bmatrix} \frac{1}{4} & 0 & -\frac{1}{4} \\ -\frac{1}{4} & 1 & -\frac{3}{4} \\ -\frac{1}{12} & 0 & \frac{5}{12} \end{bmatrix} = \begin{bmatrix} -\frac{3}{2} & 0 & -\frac{1}{2} \end{bmatrix}$$

$$\tilde{z}^* = z^* + \pi^{*T} \Delta b$$
$$= -54 + \left[-\frac{3}{2} \quad 0 \quad -\frac{1}{2} \right] \begin{bmatrix} \Delta b_1 \\ \Delta b_2 \\ \Delta b_3 \end{bmatrix} = -54 - \frac{3}{2} \Delta b_1 + 0 \Delta b_2 - \frac{1}{2} \Delta b_3$$

Let
$$\Delta b_2 = \Delta b_3 = 0$$
.
 $\Delta b_1 > 0 \Longrightarrow -\frac{3}{2} \Delta b_1 < 0 \Longrightarrow \tilde{z}^* < z^*$



$$\pi^{*T} = c_{B^*}^{T} B^{*-1}$$

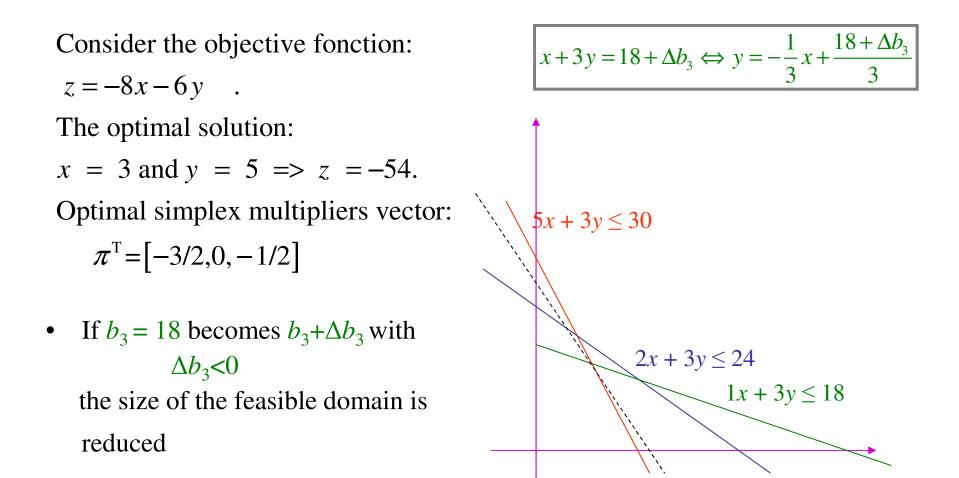
$$\pi^{*T} = \begin{bmatrix} -8 & 0 & -6 \end{bmatrix} \begin{bmatrix} \frac{1}{4} & 0 & -\frac{1}{4} \\ -\frac{1}{4} & 1 & -\frac{3}{4} \\ -\frac{1}{12} & 0 & \frac{5}{12} \end{bmatrix} = \begin{bmatrix} -\frac{3}{2} & 0 & -\frac{1}{2} \end{bmatrix}$$

$$\tilde{z}^{*} = z^{*} + \pi^{*T} \Delta b$$

= $-54 + \left[-\frac{3}{2} \ 0 \ -\frac{1}{2} \right] \begin{bmatrix} \Delta b_{1} \\ \Delta b_{2} \\ \Delta b_{3} \end{bmatrix} = -54 - \frac{3}{2} \Delta b_{1} + 0 \Delta b_{2} - \frac{1}{2} \Delta b_{3}$

Let
$$\Delta b_1 = \Delta b_2 = 0$$

 $\Delta b_3 < 0 \Longrightarrow -\frac{1}{2} \Delta b_3 > 0 \Longrightarrow \tilde{z}^* > z^*$

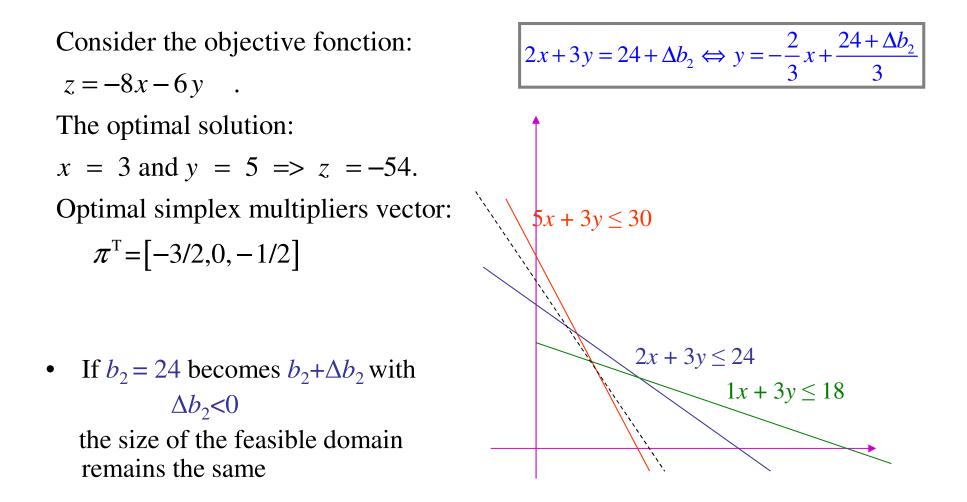


$$\tilde{z}^{*} = z^{*} + \pi^{*T} \Delta b$$

= $-54 + \left[-\frac{3}{2} \ 0 \ -\frac{1}{2} \right] \begin{bmatrix} \Delta b_{1} \\ \Delta b_{2} \\ \Delta b_{3} \end{bmatrix} = -54 - \frac{3}{2} \Delta b_{1} + 0 \Delta b_{2} - \frac{1}{2} \Delta b_{3}$

Let
$$\Delta b_1 = \Delta b_3 = 0$$

 $\Delta b_2 < 0 \Longrightarrow 0 \Delta b_2 = 0 \Longrightarrow \tilde{z}^* = z^*$



• <u>Proposition</u> In the simplex algorithm, if at some iteration, all the relative $costs \ c \ j \ge 0$ $\forall j, 1 \le j \le n$, then the current solution is optimal **Proof:** Without lost of generality, assume that the first *m* variables $x_1, x_2, ..., x_m$ are basic; i. e.,

basic var.	$x_1 \ x_2 \cdots x_r \cdots x_m \ x_{m+1} \ \cdots x_s \ \cdots x_n \ -z$	r.h.s.
<i>x</i> ₁	1 $\overline{a}_{1m+1} \cdots \overline{a}_{1s} \cdots \overline{a}_{1n}$	\overline{b}_1
x ₂	1 $\overline{a}_{2m+1}\cdots\overline{a}_{2s}\cdots\overline{a}_{2n}$	\overline{b}_2
		:
<i>x</i> ,	1 $\bar{a}_{rm+1}\cdots \bar{a}_{rs}\cdots \bar{a}_{rn}$	b,
:		:
<i>x</i> _m	$1 a_{mm+1} \cdots a_{ms} \cdots a_{mn}$	<i>b m</i>
- <i>z</i>	$\bar{c}_{m+1} \cdots \bar{c}_s \cdots \bar{c}_n = 1$	$-c_b^{\mathrm{T}}B^{-1}b$

 $\overline{z} = c_B B^{-1} b$

basic var.	x_B^T	x_R^{Γ}	- z	r.h.s.
χ_B	Ι	$B^{-1}R$	0	$B^{-1}b$
<i>Z</i>	0	$c_R^T - c_B^T B^{-1} R$	1	$-c_B^T B^{-1} b$

$$\overline{x}_i = \overline{b}_i \ge 0$$
 $i = 1, 2, ..., m$
 $\overline{x}_i = 0$ $i = m + 1, m + 2, ..., n$

basic var.	$x_1 \ x_2 \cdots x_r \cdots x_m \ x_{m+1} \ \cdots \ x_s \ \cdots \ x_n \ -z$	r.h.s.
<i>x</i> 1	1 $\overline{a}_{1m+1} \cdots \overline{a}_{1s} \cdots \overline{a}_{1n}$	$ar{b}_1$
<i>x</i> 2	1 $\overline{a}_{2m+1}\cdots\overline{a}_{2s}\cdots\overline{a}_{2n}$	\overline{b}_2
÷	· · ·	:
х,	1 $\bar{a}_{rm+1}\cdots \bar{a}_{rs}\cdots \bar{a}_{rn}$	\overline{b}_r
-	··. ·. ·.	_
x _m	$1 \overline{a}_{mm+1} \cdots \overline{a}_{ms} \cdots \overline{a}_{mn}$	b_m
- <i>z</i>	$\overline{c}_{m+1} \cdots \overline{c}_s \cdots \overline{c}_n = 1$	$-c_B^I B^{-1} b$

The objective function is as follows

$$z = 0x_1 + \ldots + 0x_m + \overline{c}_{m+1}x_{m+1} + \ldots + \overline{c}_n x_n + c_B^T B^{-1}b$$

$$\overline{z} = c_B^T B^{-1} b$$

The objective function is as follows

 $z = 0x_1 + \dots + 0x_m + \overline{c}_{m+1}x_{m+1} + \dots + \overline{c}_n x_n + c_B^T B^{-1}b$ Consider another feasible solution $\overline{x} \ge 0$ with the value $= z = c_{m+1} \overline{x}_{m+1} + c_{m+2} \overline{x}_{m+2} + \dots + c_n \overline{x}_n + c_B^T B^{-1}b$ But the hypothesis $\overline{c}_j \ge 0 \quad \forall j, 1 \le j \le n$ induces that $= z = c_{m+1} \overline{x}_{m+1} + c_{m+2} \overline{x}_{m+2} + \dots + c_n \overline{x}_n + c_B^T B^{-1}b \ge c_B^T B^{-1}b$

The objective function is as follows

$$z = 0x_1 + \dots + 0x_m + \overline{c}_{m+1}x_{m+1} + \dots + \overline{c}_n x_n + c_B^T B^{-1}b$$

Consider another feasible solution $\overline{\overline{x}} \ge 0$ with the value
$$\overline{\overline{z}} = c_{m+1}\overline{x}_{m+1} + c_{m+2}\overline{x}_{m+2} + \dots + c_n\overline{x}_n + c_B^T B^{-1}b$$

But the hypothesis $\overline{c}_j \ge 0 \quad \forall j, 1 \le j \le n$ induces that
$$\overline{\overline{z}} = c_{m+1}\overline{x}_{m+1} + c_{m+2}\overline{x}_{m+2} + \dots + c_n\overline{x}_n + c_B^T B^{-1}b \ge c_B^T B^{-1}b = \overline{z}$$

$$\ge 0$$

Then the current solution is optimal. \Box