

## Lecture 21 - catalyst

Tuesday, March 31, 2020 14:27

today : finish prox example  
 - catalyst → accelerate  
 non-convex opt.  
 - submodular opt.

### finish prox example

→ to be useful, need  $\text{prox}_\gamma^\Omega(z)$  to be efficiently computable

$$\text{prox}_\gamma^{\| \cdot \|_1}(z) = \arg \min_w \|w\|_1 + \frac{1}{2\gamma} \|w - z\|^2$$

$$\begin{aligned} \text{"soft-thresholding"} &\triangleq \begin{cases} \text{sgn}(z_i)[|z_i| - \gamma] & \text{if } |z_i| \geq \gamma \\ 0 & \text{o.w.} \end{cases} \\ &\text{(component-wise)} \end{aligned}$$

used e.g. for lasso :  $\ell_1$ -reg. least-square

FISTA → accelerated prox. gradient method

↳ SOTA for batch lasso

\* scikit-learn → uses SAGA for lasso  $\left\{ \begin{array}{l} \text{lt-reg} \\ \text{log. reg.} \end{array} \right.$

$$\text{prox SAGA} : w_{t+1} = \text{prox}_\gamma^\Omega \left( w_t - \gamma [\nabla f_{i_t}(w_t) - g_{i_t}^t + \frac{1}{n} \sum_{j=1}^n g_j^t] \right)$$

Could accelerate using "Catalyst"

Catalyst algorithm [Lin, Mairal & Harchaoui NIPS 2015]

"meta-algorithm": outer loop which uses a linearly convergent alg. in inner loop  
 to get overall acceleration (P)

main idea: use the accelerated proximal point algorithm

with approximation inner loop of prox operator

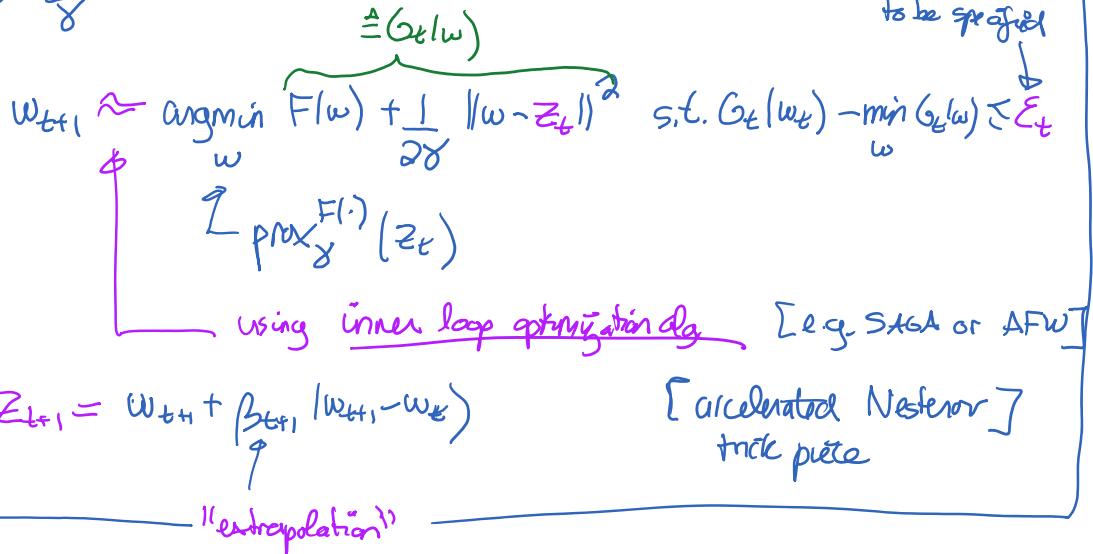
proximal pt. alg.: is proximal gradient with  $f=0$

$$w_{t+1} = \boxed{\text{prox}_\gamma^\Omega(w_t)} \quad (\text{to solve } \min_w \Omega(w))$$

Catalyst alg. (for  $\mu$ -strongly convex  $F(w)$ )

Set  $q \triangleq \frac{\mu}{\mu + \frac{1}{\gamma}}$  ( $\gamma$  is algorithmic parameter)

repeat:



$\beta_{t+1}$  is found using fancy equations so that everything works

- solve for  $\alpha_{t+1}$  in eq.  $\alpha_{t+1}^2 = (1 - \alpha_{t+1})\beta_t^2 + q\alpha_{t+1}$

$$\beta_{t+1} \triangleq \frac{\alpha_t(1 - \alpha_t)}{\alpha_t^2 + \alpha_{t+1}} \quad (\text{pick } \alpha_{t+1} \in [0, 1])$$

Catalyst trick: use  $\gamma \notin \varepsilon_t$   
 s.t. overall # of inner loop calls  
 give overall acceleration

with clever analysis of weak strong

acceleration results:

If inner loop alg. has convergence  $\exp(-\frac{\tilde{\mu}}{L}t)$   $\tilde{\mu} \geq \mu + \frac{1}{\gamma}$  (strong convexity  $G_t(w)$ )

then with correct constants

( $\mu$ -strongly convex  $F$ ) linear rate:  $\rho = \frac{1}{K} \xrightarrow{\text{becomes}} \tilde{\rho} = \frac{1}{\sqrt{K}}$  for catalyst

( $F$  convex)  $\frac{1}{t}$  on  $F$   $\xrightarrow{\text{becomes}} \frac{1}{t^2}$

results: can get (theory) accelerated

- " SAGA
- " SVRG
- " AFW
- etc...



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" AFW  
etc.. "



## non-convex optimization

recall: FW with line search on  $f$  non-convex  
 $\min_{w \in \mathbb{R}^d} g(w) \leq O\left(\frac{1}{\sqrt{\epsilon}}\right)$   
 FW gap

$$\text{convex: } \mathbb{E}[f(w_t)] - f^* \leq \epsilon$$

$$\text{non-convex: } \sqrt{\mathbb{E}\|f(w_t)\|^2} \leq \epsilon$$

gradient method:  $f(w) \leq f(w_t) + \nabla f(w_t)^T (w - w_t) + \frac{L}{2} \|w - w_t\|^2$  for  $w$

$$w_{t+1} = w_t - \frac{1}{L} \nabla f(w_t)$$

$$\Rightarrow f(w_{t+1}) \leq f(w_t) - \frac{1}{2L} \|\nabla f(w_t)\|^2$$

NIPS 2016 tutorial "Large-Scale Optimization: Beyond Stochastic Gradient Descent and Convexity"  
[Suvri Sra slides](#)

## Faster nonconvex optimization via VR

(Reddi, Hefny, Sra, Poczos, Smola, 2016; Reddi et al., 2016)

Algorithm	Nonconvex (Lipschitz smooth)
SGD	$O\left(\frac{1}{\epsilon^2}\right)$
GD	$O\left(\frac{n}{\epsilon}\right)$
SVRG	$O\left(n + \frac{n^{2/3}}{\epsilon}\right)$
SAGA	$O\left(n + \frac{n^{2/3}}{\epsilon}\right)$
MSVRG	$O\left(\min\left(\frac{1}{\epsilon^2}, \frac{n^{2/3}}{\epsilon}\right)\right)$

$$\mathbb{E}[\|\nabla g(\theta_t)\|^2] \leq \epsilon$$

### Remarks

New results for convex case too; additional nonconvex results  
 For related results, see also (Allen-Zhu, Hazan, 2016)

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## Linear rates for nonconvex problems

$$\min_{\theta \in \mathbb{R}^d} g(\theta) = \frac{1}{n} \sum_{i=1}^n f_i(\theta)$$

The Polyak-Łojasiewicz (PL) class of functions

$$g(\theta) - g(\theta^*) \leq \frac{1}{2\mu} \|\nabla g(\theta)\|^2$$

(Polyak, 1963); (Łojasiewicz, 1963)

## Linear rates for nonconvex problems

$$g(\theta) - g(\theta^*) \leq \frac{1}{2\mu} \|\nabla g(\theta)\|^2 \quad | \quad \mathbb{E}[g(\theta_t) - g^*] \leq \epsilon \quad 😎$$

Algorithm	Nonconvex	Nonconvex-PL
SGD	$O\left(\frac{1}{\epsilon^2}\right)$	$O\left(\frac{1}{\epsilon^2}\right)$
GD	$O\left(\frac{n}{\epsilon}\right)$	$O\left(\frac{n}{2\mu} \log \frac{1}{\epsilon}\right)$
SVRG	$O\left(n + \frac{n^{2/3}}{\epsilon}\right)$	$O\left((n + \frac{n^{2/3}}{2\mu}) \log \frac{1}{\epsilon}\right)$
SAGA	$O\left(n + \frac{n^{2/3}}{\epsilon}\right)$	$O\left((n + \frac{n^{2/3}}{2\mu}) \log \frac{1}{\epsilon}\right)$
MSVRG	$O\left(\min\left(\frac{1}{\epsilon^2}, \frac{n^{2/3}}{\epsilon}\right)\right)$	—

Variant of **nc-SVRG** attains this fast convergence!

(Reddi, Hefny, Sra, Poczos, Smola, 2016; Reddi et al., 2016) 22

## Submodular optimization

Submodularity is an analog of convexity for tractability

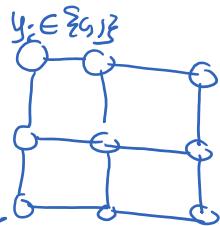
for set functions (combinatorial opt.)

$$F: 2^V \rightarrow \mathbb{R}$$

Convention here:  $F(\emptyset) = 0$

$V = \{1, \dots, d\}$  is "ground set"

$2^V = \{V \subseteq \{1, \dots, d\}\} = \text{set of all subsets of } V$



Concrete example:

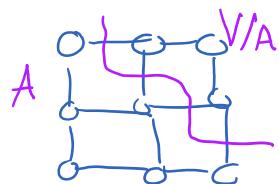
Ising model  $y_i \in \{0, 1\}$   $E(y) = \sum_i g_i y_i - \sum_{i,j \text{ neighbors}} g_{ij} y_i y_j$

when  $g_{ij} > 0$ ,  $E(y)$  is submodular  
"attractive potential"

MRF here is called "associative Markov network"  
AMN

$$F(A_y) \text{ where } A_y = \{i : y_i = 1\}$$

Can minimize  $E(y)$  by using "graph cut" alg  
(or  $F(A_y)$ )



$F$  is submodular  $\Leftrightarrow F(A) + F(B) \geq F(A \cap B) + F(A \cup B)$   $\forall A, B$

$\Leftrightarrow$  function  $A \mapsto F(A \cup \{k\}) - F(A)$  is non-increasing for all  $k$

i.e.  $F(A \cup \{k\}) - F(A) \leq F(B \cup \{k\}) - F(B)$

$\xrightarrow{\text{if } B \subseteq A}$

"diminishing return property"

$\Rightarrow$  intuitively, that greedy alg. are not "too bad" for maximization

\*  $F(A) \triangleq g(\frac{|A|}{q})$  if  $g$  is concave then  $F$  is submodular  
continuously

\* link with convexity  $\rightarrow$  Lovasz extension (cts. fct.)

\* embed sets as corners of hypercube in dimension  $d$   $v(A) = \mathbf{1}_A \in \{0,1\}^d$

Lovasz extension  $f$  extends  $F(A)$  from corners to entire hypercube  
using convex interpolation

(piecewise linear function on  $[0,1]^d$ )

$$f(w) = F(A) \text{ when } w = v(A)$$

$F$  is submodular  $\Leftrightarrow$  Lovasz extension is convex

\* can write  $f(w) = \max_{\substack{s \in \mathcal{B}(F) \\ \not\propto}} \langle s, w \rangle$   $\leftarrow$  this can be computed efficiently using greedy alg.  
"Base polytope"

$$\min_{A \in V} F(A) = \min_{w \in [0,1]^d} \left( \max_{\substack{s \in \mathcal{B}(F) \\ \not\propto}} \langle s, w \rangle \right) \rightarrow \text{use projected subgradient method}$$

$$\partial f(w) = \arg \max_{s \in \mathcal{B}(F)} \langle s, w \rangle$$

\* with  $l_2$ -regularization, use duality to get a

$$\min_{s \in \mathcal{B}(F)} \frac{1}{2} \|s\|^2$$

$\rightarrow$  use "mix-norm pt." alg.  $\odot$  SOTA for  
variant FCFW alg.  
submodular opt.