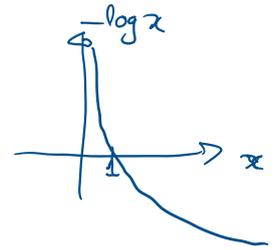
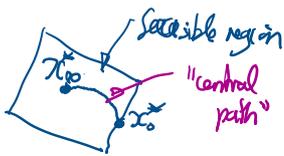


- today:
- IPM
  - duality
  - SVM dual



Interior point method:

$$\begin{aligned} \min f(x) \\ \text{s.t. } Ax=b \\ x \geq 0 \end{aligned}$$



barrier method

$$\min f(x) + t \underbrace{\left( -\sum_i \log x_i \right)}_{\text{barrier fct.}} \triangleq g_t(x)$$

s.t.  $Ax=b$

$x_t^* \triangleq \arg \min_x g_t(x)$  s.t.  $Ax=b$

$$\text{dom}(g_t(x)) = \{x : x > 0\}$$

interior point method: run Newton's method with warm start on  $g_t(x)$  and decrease  $t$  with a geometric schedule

convergence of Newton's: Newton's method take  $\log(\log(\frac{1}{\epsilon})) + \text{cst.} \cdot (g_t(x_0) - g_t(x_t^*))$  iterations  
 "self-concordant analysis"  
 to reach  $\epsilon$ -accuracy there is no condition # of  $g_t$  appearing

IPM: total # of Newton iterations is  $\log(\frac{1}{\epsilon}) \cdot \text{some constants}$   
 $\hookrightarrow$  dimensionality is appearing linearly but not the condition # of  $f$   
 (vs. gradient methods)

Review of Lagrangian duality

primal problem

$$p^* \triangleq \inf_x f(x) \\ \text{s.t. } h_i(x) \leq 0 \quad \forall i \in I$$

dual problem

$$d^* \triangleq \sup_{\alpha \geq 0} d(\alpha)$$

Sagrangian dual fct.

$$d(\alpha) \triangleq \inf_x f(x) + \sum_i \alpha_i h_i(x)$$

Fenchel duality

First: Fenchel conjugate of  $f$

$$f^*(y) \triangleq \sup_x \langle y, x \rangle - f(x)$$

$\uparrow$   
 $x^*(y)$  is s.t.  $\nabla f(x^*) = y$

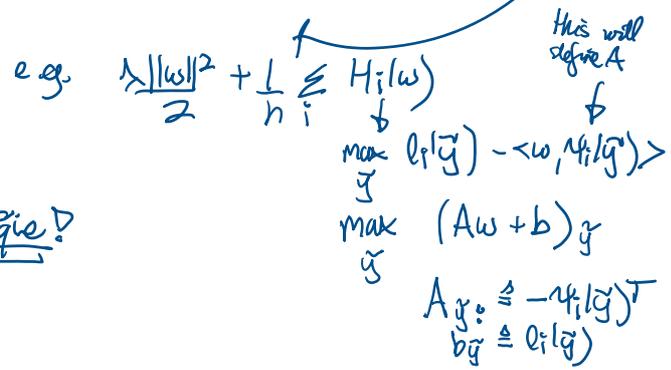
Fenchel duality:

$$p^* = \inf_{x \in \mathbb{R}^d} f(x) + g(Ax) \quad A: \mathbb{R}^d \rightarrow \mathbb{R}^p$$

$$d^* = \sup_{y \in \mathbb{R}^p} -f^*(A^T y) - g^*(y)$$

$x$

$$d^* = \sup_{y \in \mathbb{R}^p} -f^*(A^T y) - g^*(y)$$



note: not unique!

Saddle pt. perspective of Lag. duality:

$$\inf_x f(x) \text{ s.t. } h_i(x) \leq 0 \forall i \in I = \inf_x \left[ \sup_{\alpha \geq 0} \underbrace{f(x) + \sum_i \alpha_i h_i(x)}_{\substack{\text{Lagrange multiplier} \\ \mathcal{L}(x, \alpha) \text{ "Lagrangian fct."}}} \right]$$

because  $g(x) \triangleq \sup_{\alpha \geq 0} \sum_i \alpha_i h_i(x) = \begin{cases} +\infty & \text{if } x \text{ is not feasible} \\ & (\text{i.e. } \exists i \text{ s.t. } h_i(x) > 0) \\ 0 & \text{o.w.} \end{cases}$

$$p^* = \inf_x f(x) + \delta_C(x) = \inf_x \sup_{\alpha \geq 0} \mathcal{L}(x, \alpha) \geq \sup_{\alpha \geq 0} \inf_x \mathcal{L}(x, \alpha) \triangleq d^*$$

"weak duality"

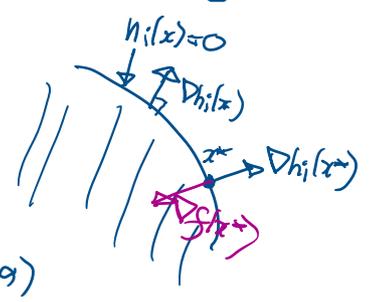
$\mathcal{L}(x, \alpha)$  "Lagrangian dual fct."

$$= \delta_C(x) \text{ where } C \triangleq \{x : h_i(x) \leq 0 \forall i\}$$

strong duality  $p^* = d^*$  (have a "saddle pt.")

KKT conditions:  $\rightarrow$  necessary conditions for  $x^*$  &  $\alpha^*$  to be primal & dual optimal (respectively) when  $f$  &  $h_i$  are differentiable fct.

- $x^*$  is primal feasible (i.e.  $h_i(x^*) \leq 0 \forall i$ )
- $\alpha^*$  is dual " (i.e.  $\alpha^* \geq 0$ )
- $\nabla f(x^*) + \sum_i \alpha_i^* \nabla h_i(x^*) = 0$  [i.e.  $(x^*, \alpha^*)$  is a stationary pt. of  $\mathcal{L}(x, \alpha)$  with respect to  $x$ ]
- $\alpha_i^* h_i(x^*) = 0 \forall i$  "complementary slackness"  
 $\rightarrow \alpha_i^* > 0$  only on active constraints  $h_i(x^*) = 0$



\*if primal problem is convex and "Slater's condition"  $\Rightarrow$  strong duality i.e.  $p^* = d^*$

\* if primal problem is convex (i.e.  $f$  &  $h_i$ 's are convex) and "Slater's condition"  $\Rightarrow$  strong duality i.e.  $p^* = d^*$

and KKT conditions are sufficient

(e.g. suppose get  $\alpha^*$  by solving dual then get  $x^*$  through KKT conditions  $\Rightarrow x^*$  solves primal)

$\exists x$  in rel interior (dom  $f$ )  
 s.t.  $h_i(x) < 0 \forall i$   
 when  $h_i$  is not affine

$\rightarrow$  see counter-example for strong duality (see Boyd's book example 5.21)

min  $f(x, y)$   
 s.t.  $\frac{x^2}{y} \leq 0$

where  $f(x, y) = \begin{cases} x^2/y & y > 0 \\ +\infty & \text{otherwise} \end{cases}$

violates Slater's cond. and unbounded, no strong duality

other comments:

- $f(x) \geq d(\alpha) \forall x \text{ \& } \alpha$  feasible
- $d(\alpha)$  is always concave in  $\alpha$
- $d(\alpha)$  is obtained through unconstrained opt.
- no unique "Lagrangian dual problem" e.g.  $h_i(x) \leq 0 \Leftrightarrow \alpha h_i(x) \leq 0$  for some  $\alpha > 0$

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SVM struct dual problem

consider primal: min  $\frac{\Delta \|w\|^2}{2} + \frac{1}{n} \sum_{i=1}^n \xi_i$

s.t.  $\frac{1}{n} \xi_i \geq \frac{1}{n} H_i(\tilde{y}; w) \forall i \in \{1, \dots, n\}$   
 $\hookrightarrow \xi_i(\tilde{y}) - w^T \psi_i(\tilde{y})$

$L(w, \xi; \alpha) = \frac{\Delta \|w\|^2}{2} + \frac{1}{n} \sum_i \xi_i + \frac{1}{n} \sum_{i, \tilde{y}} \alpha_i(\tilde{y}) [ \xi_i(\tilde{y}) - w^T \psi_i(\tilde{y}) - \xi_i ]$

$\frac{\partial L}{\partial \xi_i} = 0 \implies \sum_{\tilde{y}} \alpha_i(\tilde{y}) = 1$  want  $h_{i, \tilde{y}}(w, \xi) \leq 0$

i.e.  $\alpha_i \in \Delta_{|\mathcal{Y}|}$

$\nabla_w L = 0 \implies \Delta w - \frac{1}{n} \sum_{i, \tilde{y}} \alpha_i(\tilde{y}) \psi_i(\tilde{y}) = 0$   
 $\rightarrow$  if  $\alpha_i \in \Delta_{|\mathcal{Y}|}; d(\alpha) = +\infty$

$$\nabla_w L = 0 \Rightarrow Aw - \sum_{i, \tilde{y}} \alpha_i(\tilde{y}) \psi_i(\tilde{y}) = 0$$

$$w(\alpha) = A\alpha \Rightarrow \boxed{w(\alpha) \triangleq \sum_{i, \tilde{y}} \alpha_i(\tilde{y}) \psi_i(\tilde{y})} \quad \text{KKT conditions}$$

(here  $w^* = w(\alpha^*)$ )

$$d(\alpha) = L(w(\alpha), \xi(\alpha); \alpha) = \frac{\lambda}{2} \|w(\alpha)\|^2 + \sum_{i, \tilde{y}} \alpha_i(\tilde{y}) \underbrace{[b_i(\tilde{y}) - w^T \psi_i(\tilde{y})]}_{\langle b, \alpha \rangle}$$

notation:  $\alpha \in \Delta_{\mathcal{Y}_1} \times \Delta_{\mathcal{Y}_2} \times \dots \times \Delta_{\mathcal{Y}_n}$

$b \in \mathbb{R}^m$   
 $\alpha \in \mathbb{R}^m$  exponentially big  
 $m = \sum_i |\mathcal{Y}_i|$

$$b = (b_i)_{i=1}^n \quad b_i(\tilde{y}) \triangleq \sum_n \ell_i(\tilde{y})$$

$$A \in \mathbb{R}^{d \times m}$$

$$A \triangleq \left[ \dots \underbrace{\sum_{i, \tilde{y}} \psi_i(\tilde{y})}_{\substack{\text{column} \\ \psi(i, \tilde{y})}} \dots \right]$$

$$\boxed{w(\alpha) = A\alpha}$$

$$d(\alpha) = \langle b, \alpha \rangle + \frac{\lambda}{2} \|w(\alpha)\|^2 - w(\alpha)^T \left( \sum_{i, \tilde{y}} \alpha_i(\tilde{y}) \psi_i(\tilde{y}) \right)$$

$\lambda w(\alpha)$   $\otimes$

$$d(\alpha) = \langle b, \alpha \rangle - \frac{\lambda}{2} \|w(\alpha)\|^2$$

dual problem:  $\max_{\substack{\alpha_i \in \Delta_{\mathcal{Y}_i} \\ 1 \leq i \leq n}} \frac{\lambda}{2} \|A\alpha\|^2 + \langle b, \alpha \rangle$

(exponentially many dual variables)

complementary slackness:  $\alpha_i^*(\tilde{y}) > 0 \Rightarrow h_{i, \tilde{y}}(w^*, \xi^*) = 0$

ie.  $\xi_i^* = b_i(\tilde{y}) - w^{*T} \psi_i(\tilde{y}) = H_i(\tilde{y}; w^*)$

define  $\xi_i(\alpha) = \max_{\tilde{y}} H_i(\tilde{y}; w(\alpha))$

by primal feasibility

$$\xi_i^* \geq H_i(\tilde{y}^*, w^*) \quad \forall \tilde{y}^* \in \mathcal{Y}_i$$

$$\Rightarrow \xi_i^* = \max_{\tilde{y} \in \mathcal{Y}_i} H_i(\tilde{y}; w^*) = H_i(w^*) = \xi_i(\alpha^*)$$

Support vectors are when  $\alpha_i^*(\tilde{y}) > 0$  for  $\tilde{y} \neq \tilde{y}^{(i)}$

[necessary condition:  $\tilde{y} \in \arg \max_{\tilde{y} \in \mathcal{Y}_i} H_i(\tilde{y}; w^*)$ ]

$\alpha^*$  is sparse

(aside:  $y^{(i)}$  has  $\psi_i(y^{(i)}) = 0$ )

→ for input examples  $i$  s.t.  $\arg\max H_i(\tilde{y}; w^*) = \{y^{(i)}\}$

then  $\alpha_i(y^{(i)}) = 1 \Rightarrow$  this input example  
does not influence  $w^*$

i.e. not a "support vector"