

today: • MCMC
• review Markov chains
• M-H. alg.

MCMC - Markov chain Monte-Carlo

idea: is to relax indep. assumption samples
to allow adaptive proposal dist.

ie. we'll run a chain $X_t | X_{t-1}$ st. $X_t \xrightarrow{t \rightarrow \infty}$ in dist. to target dist. p

'stationary dist. of a chain'

then we can approximate

$$\mathbb{E}_p[f(x)] \approx \frac{1}{T-T_0} \sum_{t=T_0+1}^T f(X_t)$$

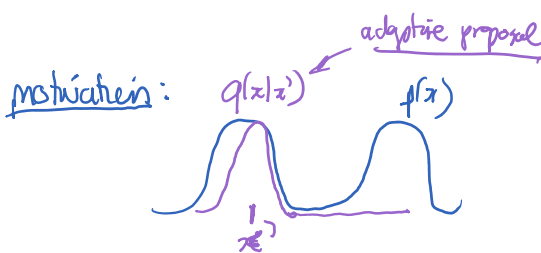
T_0 is called "burn-in period" $\rightarrow$ depends on "mixing time" of Markov chain

⊛ no need to thin the samples [ie. use Δt between samples to get more independence]

as this yield higher variance

$\rightarrow$ better to use all samples after T_0 to estimate μ

[unless it is too expensive]



before: samples were $x^{(t)} \sim q$

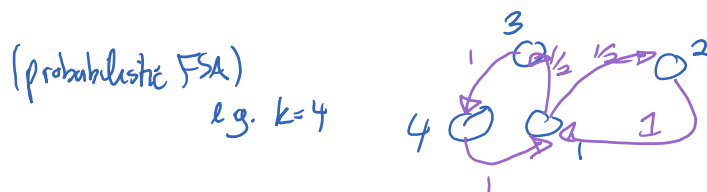
$$\text{MCMC } x^{(t)} | x^{(t-1)} \sim q(\cdot | x^{(t-1)})$$

Markov transition prob

Review of (finite state space) Markov chain [$|X|=k$]

• as DGM, $x^{(0)} \rightarrow x^{(1)} \rightarrow x^{(2)} \rightarrow \dots \rightarrow x^{(t-1)} \rightarrow x^{(t)}$

• there is also transition prob. pt. of view: use one node per state



[homogeneous M.C.]

↳ i.e. $P\{X_t = i | X_{t-1} = j\} = A_{ij}$ (no time dep.)

A is a $k \times k$ matrix s.t. $\mathbb{1}_k^T A = \mathbb{1}_k^T$

"left-stochastic matrix"

vector of ones of size k

(*) (as in HM) suppose $P\{X_{t-1} = j\} = (\pi)_j$

$$P\{X_t = i\} = \sum_j P\{X_t = i | X_{t-1} = j\} P\{X_{t-1} = j\}$$

A_{ij} π_j

$$\pi_t = A \pi_{t-1}$$

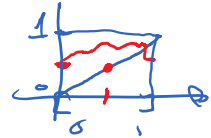
$$\Rightarrow \boxed{\pi_t = A^t \pi_0}$$

stationary dist. π of A is a dist. π s.t. $A\pi = \pi$

[note that this π is a right e-vector of A with e-value 1]

fact: every stochastic matrix has at least 1 stat. dist?

[by Brouwer's fixed pt. thm.]



def: irreducible Markov chain $\Leftrightarrow$ there exists a positive prob "path" from any i to any j (states)

$$\forall (i,j), \exists \text{ an } \overset{\text{(positive)}}{\text{integer}} m_{ij} \text{ s.t. } (A^{m_{ij}})_{ji} > 0$$

(by Perron-Frobenius thm. $\Rightarrow$ irreducible M.C. has a unique stat. dist.)

[multiplicity of e-value 1 is 1]

(*) in order to converge to it, we need aperiodicity as well

irreducible and aperiodic M.C. $\Leftrightarrow \exists$ an integer m s.t. $(A^m)_{ij} > 0$

aka regular M.C. (finite state space)

$$(A^m)_{ij} > 0 \forall i,j$$

or ergodic M.C.

(*) [note: a sufficient condition for an irreducible M.C. to be aperiodic is $\exists i$ s.t. $A_{ii} > 0$]

ω be aperiodic i.e. $\exists i$ st. $A_{ii} > 0$

example of a regular M.C. on k states

$$A = \begin{pmatrix} 0 & \frac{1}{2} & \frac{1}{2} \\ \frac{1}{2} & 0 & \frac{1}{2} \\ \frac{1}{2} & \frac{1}{2} & 0 \end{pmatrix} = \frac{1}{k-1} (11^T - I)$$

$$A^2 = \frac{1}{(k-1)^2} (11^T 11^T - 211^T + I) = \frac{1}{(k-1)^2} ((k-2)11^T + I) \text{ for } k \geq 3, \text{ this } > 0$$

[but, for $k=2$, it is not aperiodic, $A^2 = I$ $A = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$]

Thm: if a finite M.C. is ergodic (regular)

then $\exists$ a unique stat. dist. π

and for any starting dist. π_0 , $\lim_{t \rightarrow \infty} A^t \pi_0 = \pi$

the speed of convergence is related to the mixing time τ of the chain

$$\tau \triangleq \frac{1}{1 - |\lambda_2(A)|}$$

$\lambda_2(A)$ 2nd biggest λ -value of A

$$\|A^t \pi_0 - \pi\|_1 \leq C \exp(-t/\tau)$$

after τ steps, error decreases by $\frac{1}{e}$

16 Nov

⊛ intuition (from linear algebra) [informal argument]

simpler case, suppose A is symmetric & psd., by spectral thm., A is diagonalizable with orthogonal matrix U

$$A = U \Sigma U^T \text{ with } \Sigma = \begin{pmatrix} \lambda_1 & & 0 \\ & \ddots & \\ 0 & & \lambda_k \end{pmatrix}$$

[by Perron-Frobenius thm.]

$U \rightarrow$ orthonormal basis

of $\mathbb{R}^k$ vectors

$$U U^T = U^T U = I \quad U = (u_1 \dots u_k)$$

$$\lambda_1 = 1 > |\lambda_2| \geq \dots \geq |\lambda_k|$$

$$u_1 = \frac{1}{\sqrt{k}} \mathbf{1} \quad A \pi = \pi$$

let α_0 be coordinates of π_0 in U basis

$$\pi_0 = U \alpha_0 \quad (\alpha_0 = U^T \pi_0)$$

$$A^t \pi_0 = (u_1 u_1^T) (u_2 u_2^T) \dots (u_k u_k^T) (u_{k+1} u_{k+1}^T)$$

$$= u \Sigma^t \alpha_0$$

$$= (\alpha_0)_1 \frac{1}{\|u_1\|_2} u_1 + (\alpha_0)_2 \frac{1}{\|u_2\|_2} u_2 + \dots + (\alpha_0)_k \frac{1}{\|u_k\|_2} u_k$$

$$\Sigma^t = \begin{pmatrix} 1 & & & \\ & \lambda_1^t & & \\ & & \lambda_2^t & \\ & & & \ddots \\ 0 & & & & \lambda_k^t \end{pmatrix}$$

(if $\frac{(\alpha_0)_1}{\|u_1\|_2} = 1$) [fishy; not true in general]

$$\|A^t \pi - \pi\|_1 = \|(\alpha_0)_2 \lambda_2^t u_2 + \dots\| \leq C |\lambda_2|^t \quad |\lambda_2| = 1 - \epsilon_1$$

first e-gap

$$\epsilon_1 = 1 - |\lambda_2|$$

$$|\lambda_2| \leq \exp(-\epsilon_1)$$

$$|\lambda_2|^t \leq \exp(-t \epsilon_1)$$

$$\frac{1}{t} \Rightarrow \gamma = \frac{1}{1 - |\lambda_2|}$$

* mixing time is often (usually) exponentially big?

* How we design A s.t. $A^t \pi_0 \rightarrow \pi$?

one "easy way"

reversible M.C.

$\Leftrightarrow \exists \text{ dist. } \pi \text{ s.t.}$

$$A_{ij} \pi_j = A_{ji} \pi_i \quad \forall (i,j)$$

"detailed balance equation"

this is a sufficient condition (but not nec.) to get $A\pi = \pi$

it means $P\{X_{t+1}=i | X_t=j\} = \pi_i$
then $P\{X_t=i, X_{t+1}=j\} = \pi_i A_{ij}$
 $1 = P\{X_t=j, X_{t+1}=i\} = \pi_j A_{ji}$

$$\text{proof: } (A\pi)_i = \sum_j A_{ij} \pi_j \stackrel{\text{detailed balance}}{=} \sum_j A_{ji} \pi_i = \pi_i \left(\sum_j A_{ji} \right)$$

Metropolis-Hastings alg.:

goal $\rightarrow$ construct a M.C. with det. dest. $p(x)$ [our target]

[assume $p(x) > 0 \quad \forall x$]

use some proposal $q(x'|x)$

accept new state x' with prob.

if reject it $\rightarrow$ stay in same state x

[this is still new sample]

vs.

rejection sampling where only "accepted states" are new samples

no dependence on normalization of p

$$\alpha(x'|x) \triangleq \min \left\{ 1, \frac{q(x|x') p(x')}{q(x'|x) p(x)} \right\}$$

acceptance ratio to satisfy detailed balance

vs. rejection sampling where only "accepted states" are new samples

M.H. alg.

start at $x^{(0)}$

for $t=1, \dots$,
 • propose $x^{(t)} \sim q(x' | x^{(t-1)})$ important design choice
 flip a coin, with prob. $a(x^{(t)} | x^{(t-1)})$ to be 1
 • if accept (coin = 1)
 let $x^{(t)} = x^{(t)}$
 o.w.
 let $x^{(t)} = x^{(t-1)}$
 end for

note: for symmetric proposal $q(x'|x) = q(x|x')$, always accept if $p(x') \geq p(x)$

[Metropolis' alg.]

→ like noisy hill-climbing alg.

[verify as exercise that M.H. satisfies the detailed eq. with $\pi = p^{\text{target}}$]

⊛ for convergence: if M.H. chain is ergodic, then converge to unique stst. dist.

sufficient conditions

- for irreducibility $q(x'|x) > 0 \forall x' \neq x \in X$
- for aperiodicity
 - either $q(x|x) > 0$ for some $x \in X$
 - or $a(x'|x) < 1$ for some $x \in X$ and x' st. $q(x'|x) > 0$

↓
 $A_i > 0$
 for some i

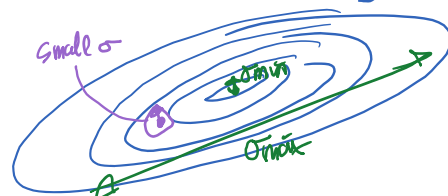
⊛ aside: it is still ok to change proposal with time
 $q_t(x'|x)$

as long as choice of q_t does not depend on $x^{(t-1)}$

then convergence theory above will go through

slow mixing example:

suppose p is a $N(\mu, \Sigma)$ & $q(x'|x) = N(x'|x, \sigma^2 I)$



"small move" → more steps needed



* high prob. of rejection



"small moves" $\Rightarrow$ many steps needed

here best mixing time related to
ratio $\frac{\sigma_{\max}}{\sigma_{\min}}$

reference for mixing times:

Markov Chains and Mixing Times

David A. Levin, Yuval Peres, Elizabeth L. Wilmer

<https://pages.uoregon.edu/dlevin/MARKOV/markovmixing.pdf>